

Authorisation in MAFTIA



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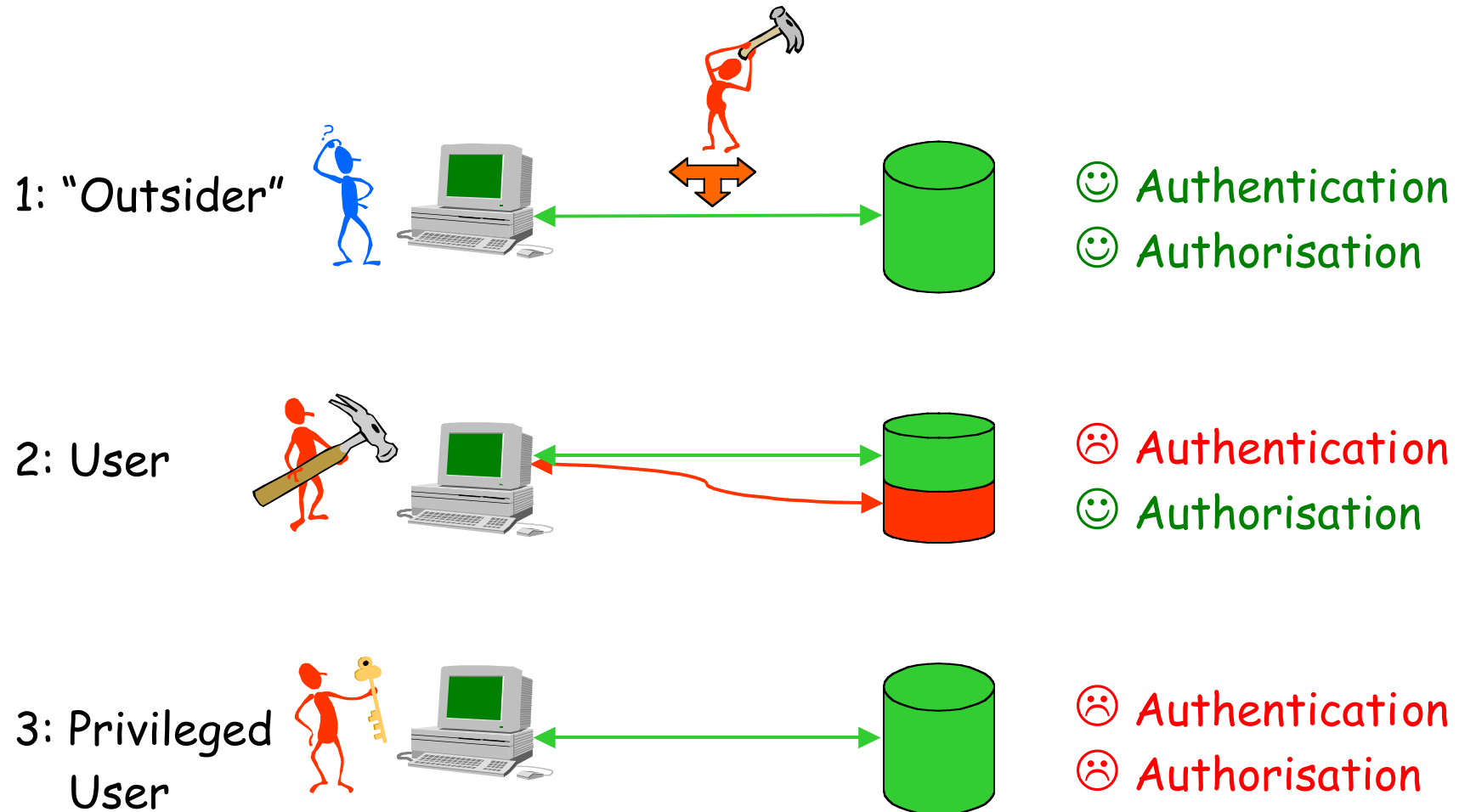


MAFTIA Workplan



- ❖ WP1: Conceptual model and architecture
- ❖ WP2: Dependable middleware
- ❖ WP3: Intrusion detection
- ❖ WP4: Dependable trusted third parties
- ❖ WP5: Distributed authorisation
- ❖ WP6: Assessment

Who are the intruders?



because the "least privilege principle" is not implemented

Intrusion Tolerance



Intrusion into a part of the system should give access only to non-significant, non-sensitive information

FRS: Fragmentation-Redundancy-Scattering

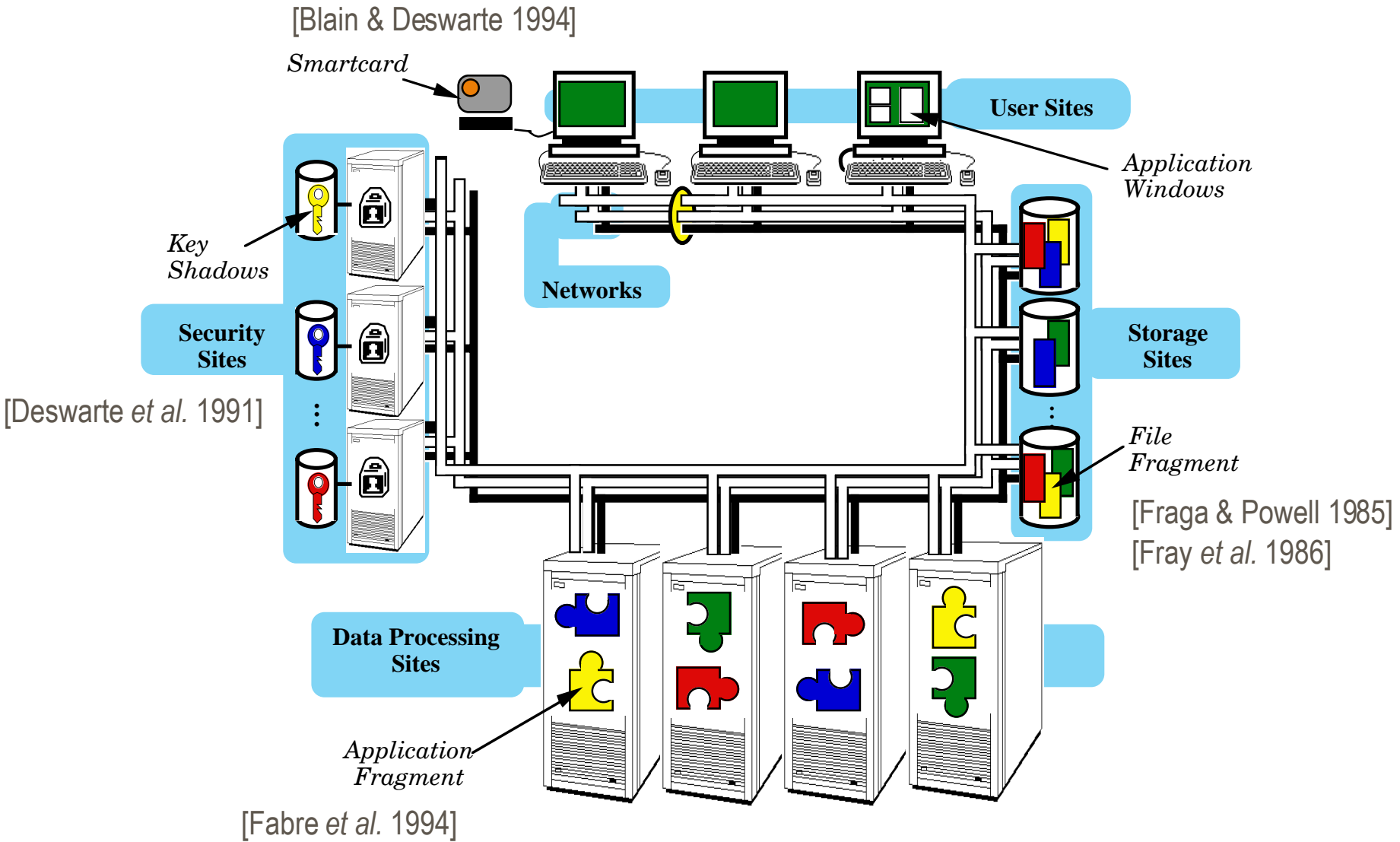
- **Fragmentation**: split the data into fragments so that isolated fragments contain no significant information: *confidentiality*
- **Redundancy**: add redundancy so that fragment modification or destruction would not impede legitimate access: *integrity + availability*
- **Scattering**: isolate individual fragments

Different kinds of scattering

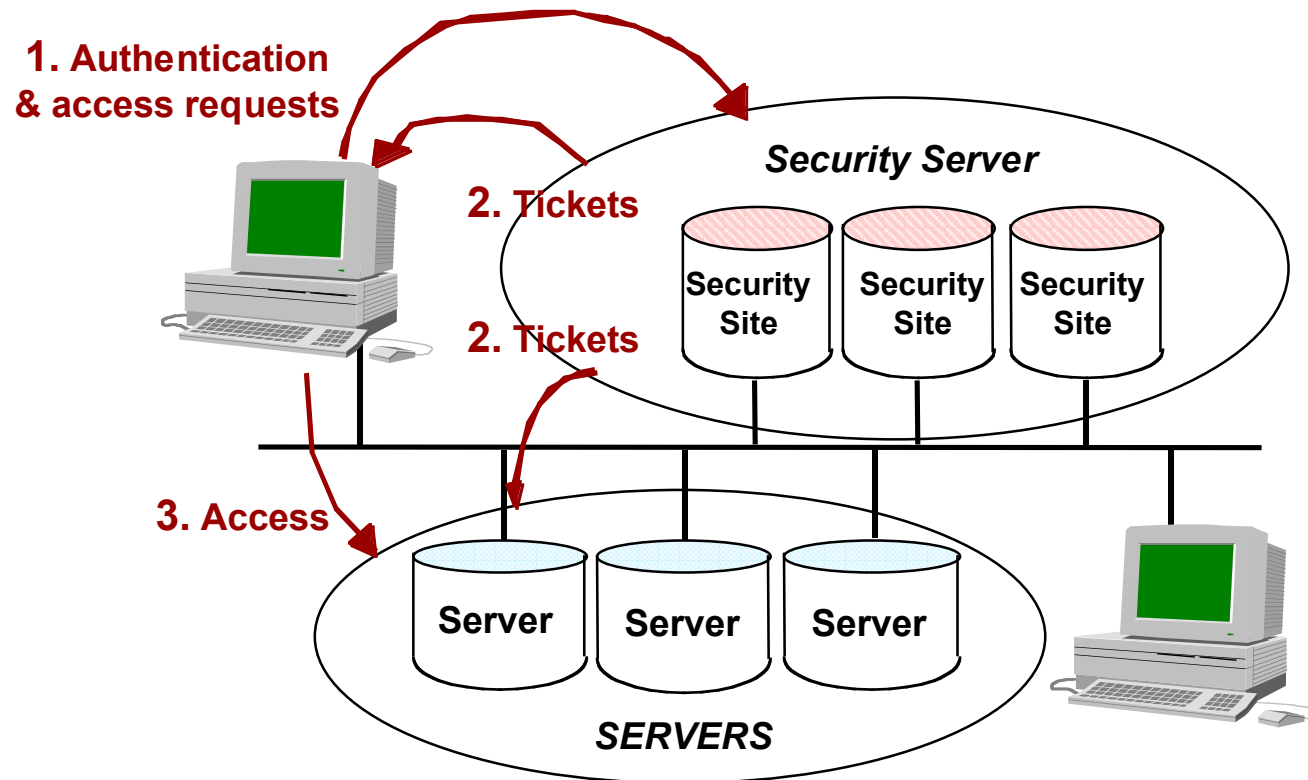


- ❖ **Space:** use different transmission links and different storage sites
- ❖ **Time:** mix fragments (from the same source, from different sources, with jamming)
- ❖ **Frequency:** use different carrier frequencies (spread-spectrum)
- ❖ **Privilege:** require the co-operation of differently privileged entities to realise an operation (separation of duty, secret sharing)

Prototype (1986-1996)

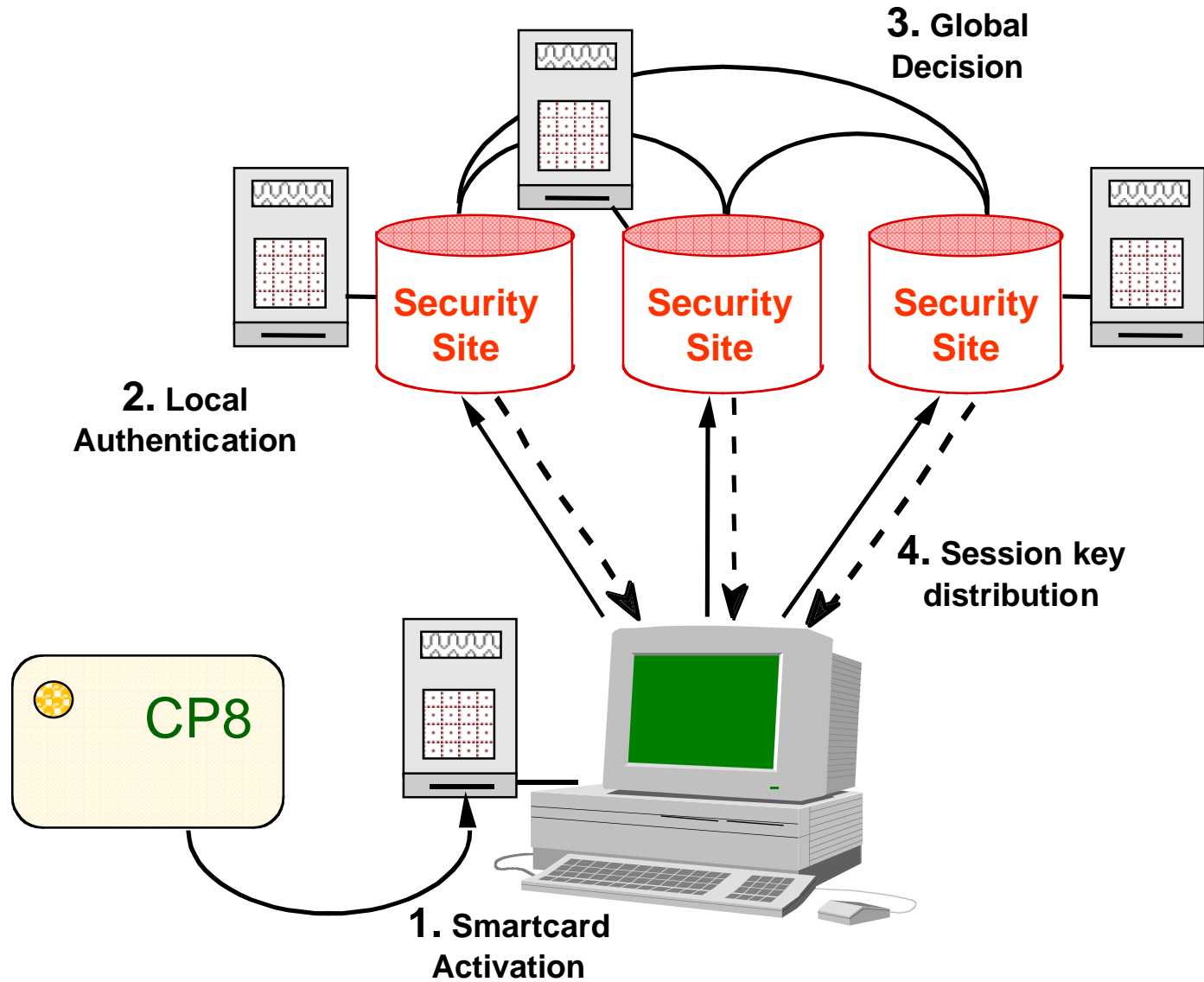


FRSed Security Management

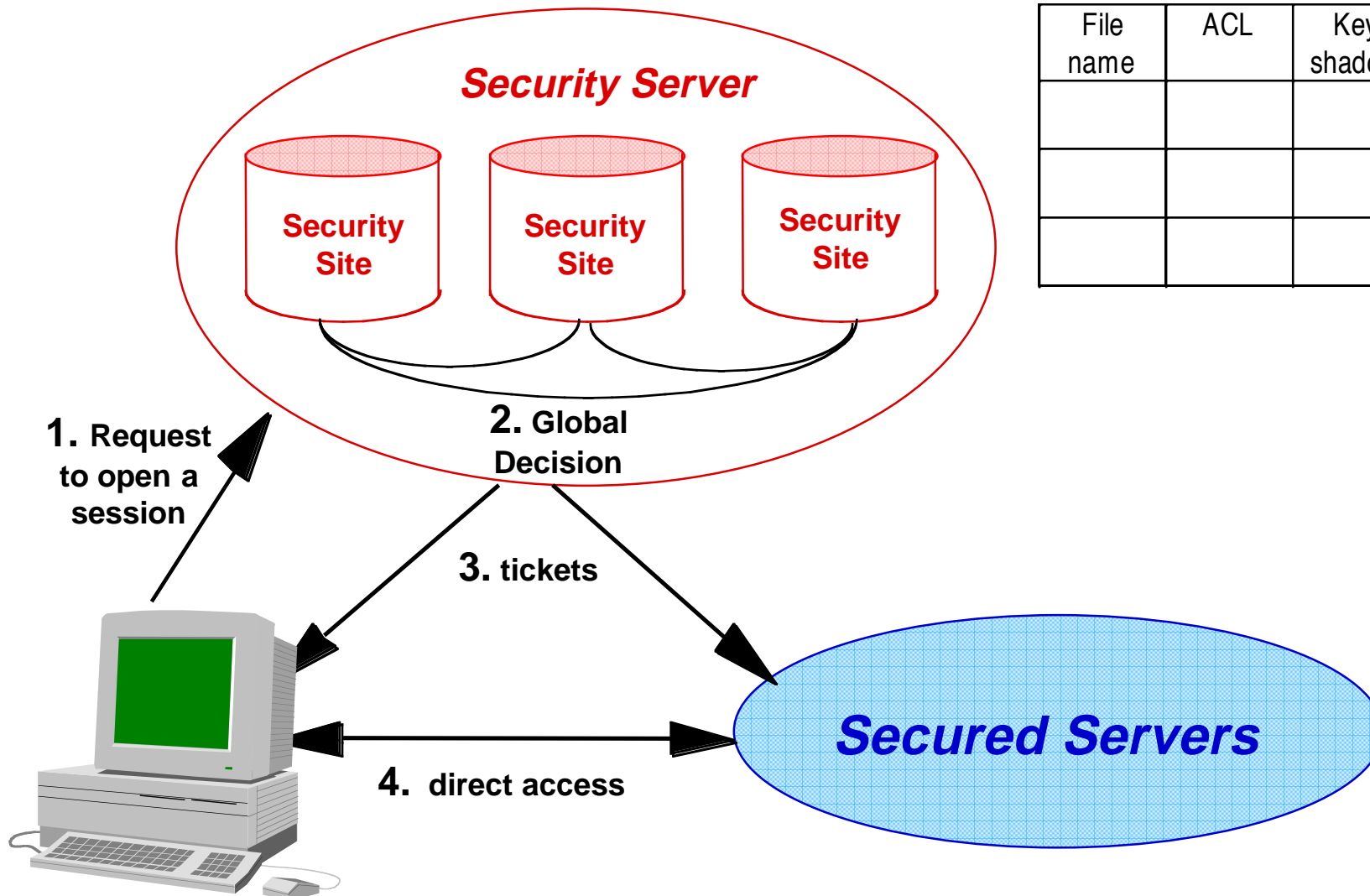


- No single trusted site or administrator
- Global trust in a majority of security sites (and administrators)

Authentication

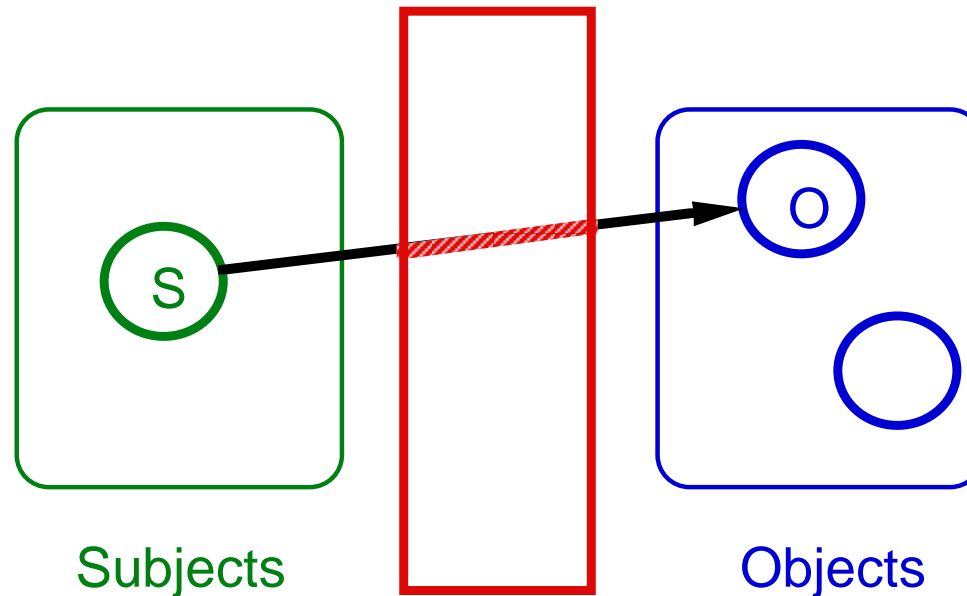


Authorisation



File name	ACL	Key shadow

Authorisation: reference monitor



Subjects

Reference
Monitor

Objects

Subject (active)
≈ process

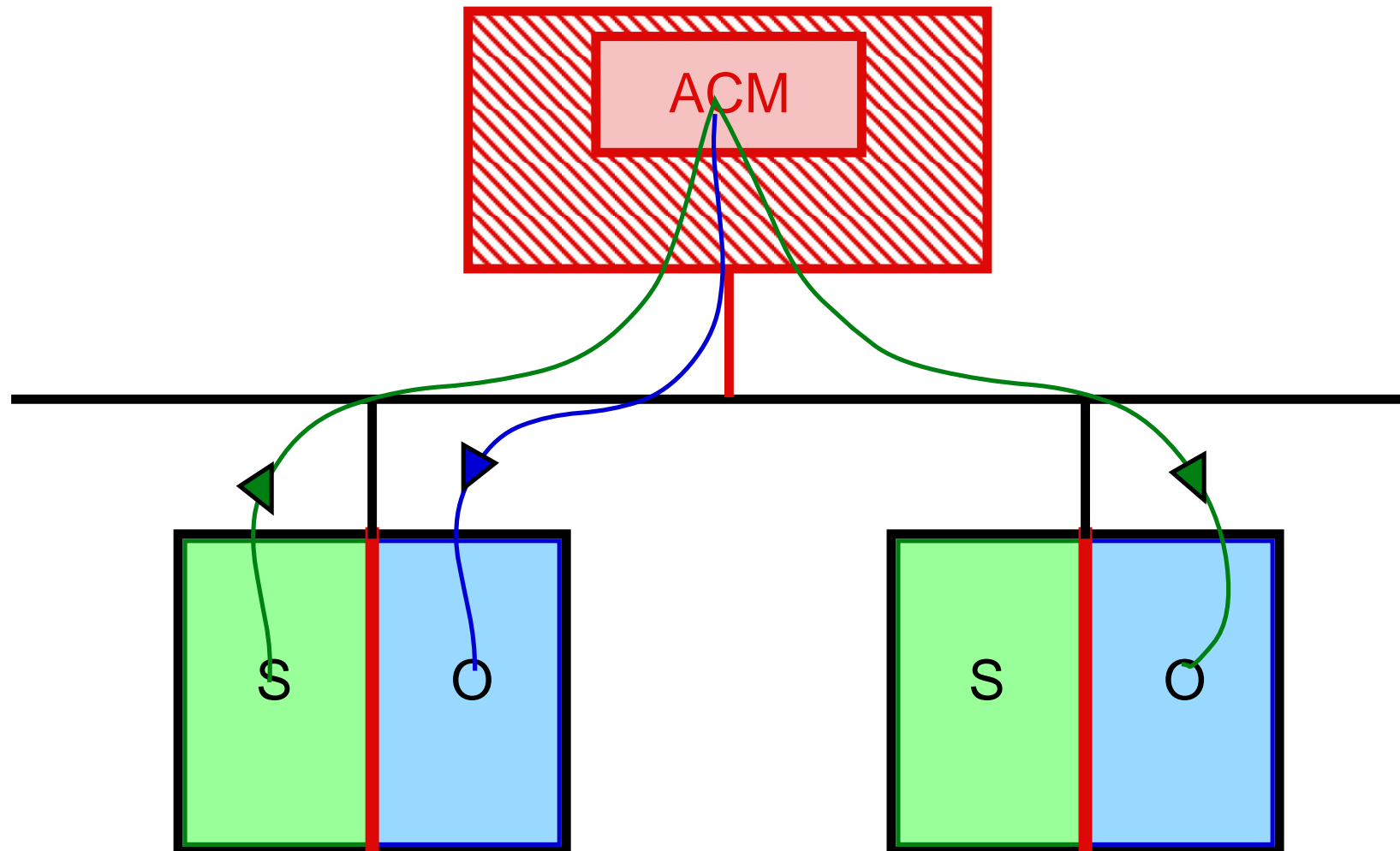
Object (passive)
≈ information containers

- Reference Monitor:
- tamperproof
 - always invoked
 - verifiable as correct

Distributed Authorisation ? (1)



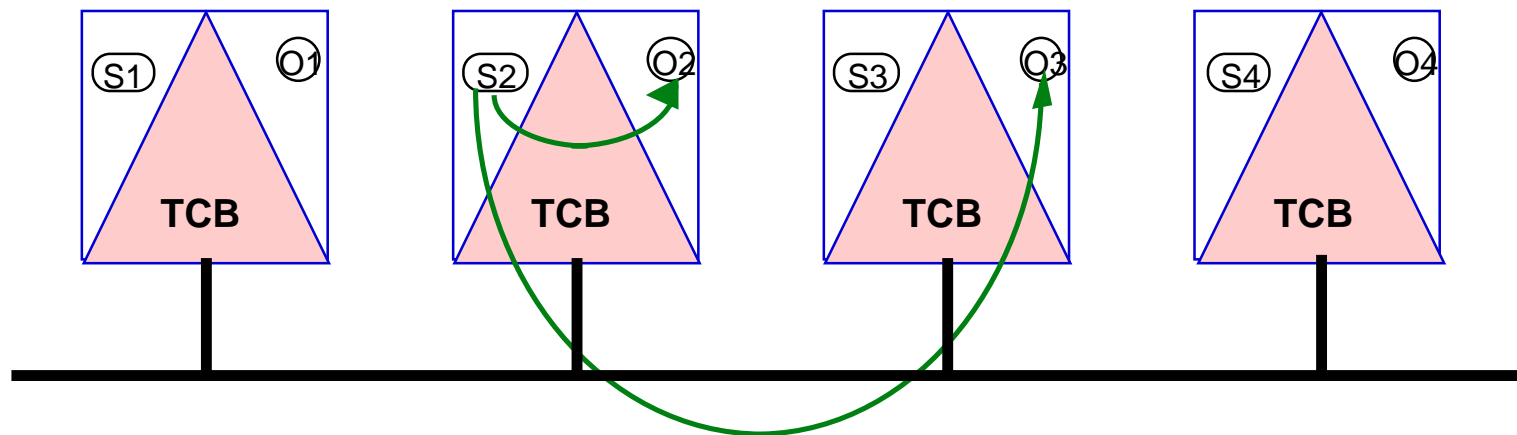
Reference Monitor



Distributed Authorisation ? (2)



Red Book (TNI)



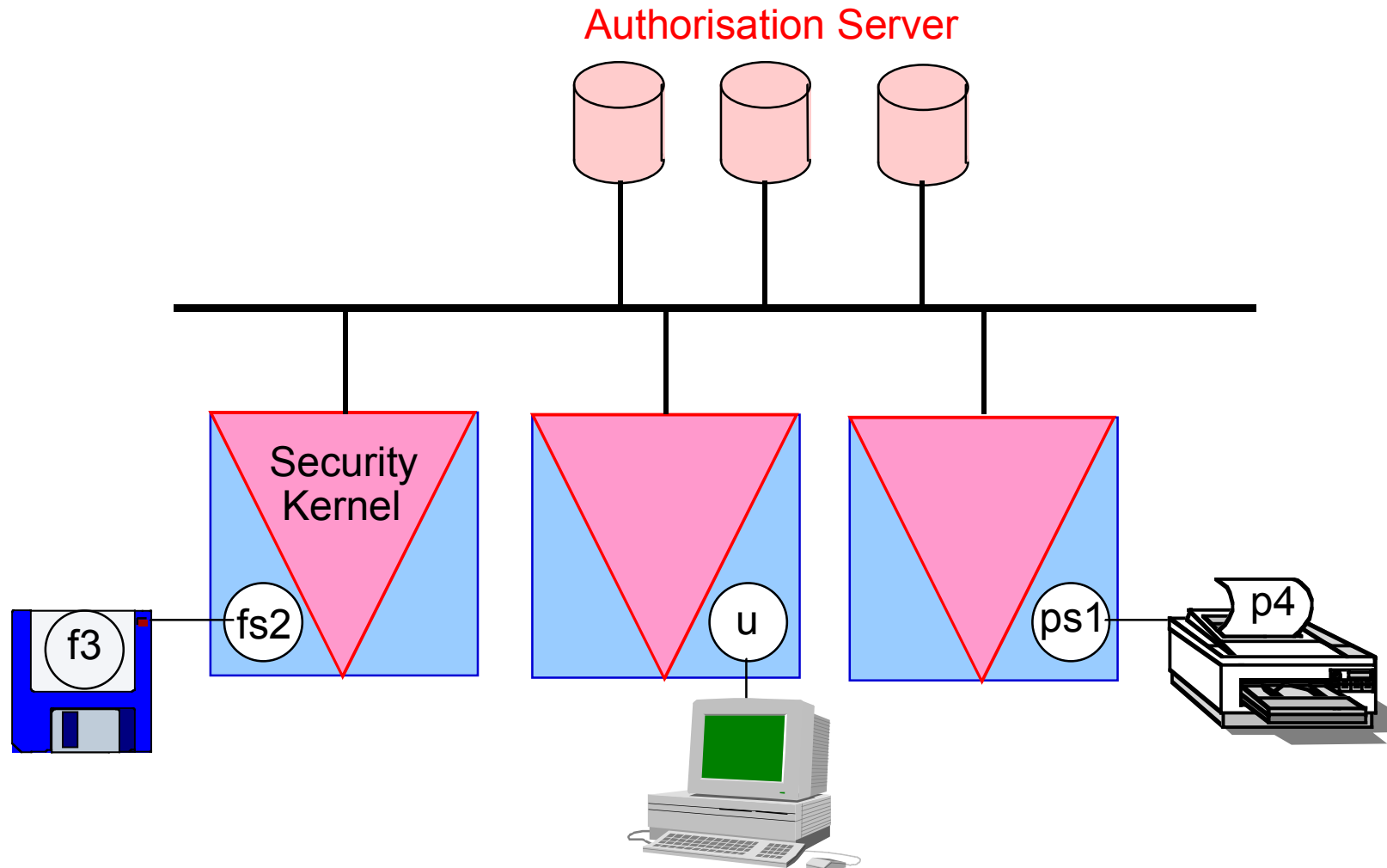
☺ No bottleneck, no single-point-of-failure

☹ Mutual trust between TCBs, consistency?

Authorisation Scheme for DOOS



[Nicomette & Deswarte 1997]



Authorisation Scheme



Access Matrix :

Method rights: corresponding to the authorisation for an object to call another object methods

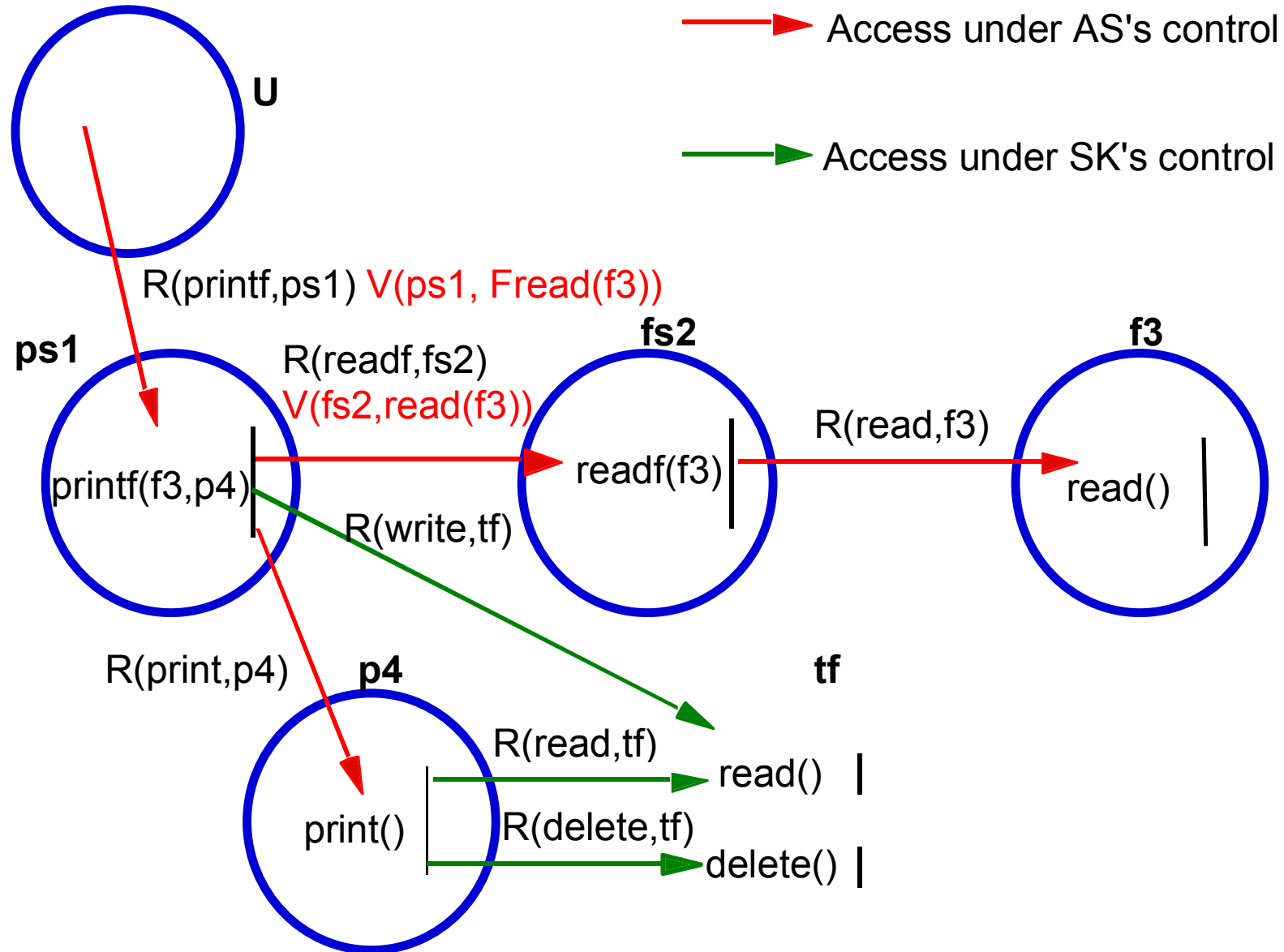
Symbolic rights: corresponding to the authorisation for an object to execute high level operations

	ps1	fs2	f 3	p 4
u			<code>PF(this, PRINTER)</code>	<code>PF(FILE, this)</code>
ps1				print
fs2				

Symbolic right rules: to check authorisation for high level operations

Capability creation rules: to grant capabilities and *vouchers* to enable high level operations

Example: $u :: PF(f3, P4)$



MAFTIA Authorisation



- ❖ Intrusion tolerant authorisation servers
 - ❖ Multi-party transactions
(not simple client-server relations)
 - ❖ Local protection

IT Authorisation Servers



Same technique as for FRS Security servers:

- ❖ Non-confidential information is replicated
- ❖ Confidential information is fragmented (threshold crypto)
- ❖ Global consensus (majority voting or Byzantine agreement)
- ❖ Distribution of capabilities/vouchers (threshold crypto)

Multi-party transactions



Based on DOOS Authorisation Scheme,
but...

- ❖ Implementing "separation of duty"
(transactions initiated jointly by several
users)

Local protection



- ❖ Internet applications => no modification of user workstations
- ❖ No security kernel, but JVM (?)
- ❖ Capabilities/vouchers -> applets
- ❖ Possibly enforced by smartcards (e.g., JavaCards) ... which are already necessary for authentication

Use cases / Scenarios



❖ Medical:

French Healthcare Professional Network

- Patient smartcard Vitale
- Professional smartcard CPS
- Electronic prescriptions
- Network, PKI, anonymisation, delegation, ...

❖ Electronic Commerce:

Auction services, privacy enforced payment, ...

References



- ❖ Blain, L. and Deswarte, Y. (1994). A Smartcard Fault-Tolerant Authentication Server, in *1st Smart Card Research and Advanced Application Conference (CARDIS'94)*, Lille, France, pp.149-165.
- ❖ Deswarte, Y., Blain, L. and Fabre, J.-C. (1991). Intrusion Tolerance in Distributed Systems, in *Symp. on Research in Security and Privacy*, Oakland, CA, USA, pp.110-121.
- ❖ Deswarte, Y., Fabre, J.-C., Laprie, J.-C. and Powell, D. (1986). A Saturation Network to Tolerate Faults and Intrusions, in *5th Symp. on Reliability of Distributed Software and Database Systems*, Los Angeles, CA, USA, pp.74-81, IEEE Computer Society Press.
- ❖ Fabre, J.-C., Deswarte, Y. and Randell, B. (1994). Designing Secure and Reliable Applications using FRS: an Object-Oriented Approach, in *1st European Dependable Computing Conference (EDCC-1)*, Berlin, Germany LNCS 852, pp.21-38.
- ❖ Fraga, J. and Powell, D. (1985). A Fault and Intrusion-Tolerant File System, in *IFIP 3rd Int. Conf. on Computer Security*, (J. B. Grimson and H.-J. Kugler, Eds.), Dublin, Ireland, Computer Security, pp.203-218.
- ❖ Fray, J.-M., Deswarte, Y. and Powell, D. (1986). Intrusion-Tolerance using Fine-Grain Fragmentation-Scattering, in *Symp. on Security and Privacy*, Oakland, CA, USA, pp.194-201.
- ❖ Nicomette, V. and Deswarte, Y. , An Authorization Scheme for Distributed Object Systems, in *Symp. on Security and Privacy*, Oakland, CA, USA, pp. 21-30.

Commercials



- ❖ MAFTIA: <http://www.research.ec.org/maftia/>
- ❖ SQUALE: <http://www.research.ec.org/squale/>
- ❖ Int. Conf. Dependable Systems and Networks (DSN 2000): 26-28 June 2000, New York
<http://www.dependability.org>
 - Workshop on Dependability despite malicious faults
- ❖ Recent Advances in Intrusion Detection (RAID 2000): 2-4 October 2000, Toulouse
<http://www.raid-symposium.org>
- ❖ European Symp. on Research in Computer Security (ESORICS 2000): 4-6 October 2000, Toulouse
<http://www.esorics.org>