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Modeling Microprocessor Faults on High-Level Decision Diagrams

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Outline

- **Introduction**
- **Motivations and contributions**
- **Discussion: faults and tests**
- **Fault modeling with Decision Diagrams**
- **Modeling microprocessor faults**
- **Experimental results**
- **Conclusions**

Introduction

- **Fault models are needed for**
 - test generation,
 - test quality evaluation and
 - fault diagnosis
- **To handle real physical defects is too difficult**
- **The fault model should**
 - reflect accurately the behaviour of defects, and
 - be **computationally efficient**
- **Usually combination of different fault models is used**
- **Fault model free approaches (!)**

Introduction

- **Fault modeling levels**

- Transistor level faults
- Logic level faults
 - stuck-at fault model
 - bridging fault model
 - open fault model
 - delay fault model
- Register transfer level faults
- ISA level faults (MP faults)
- SW level faults

Low-Level models

High-Level models

- Hierarchical fault handling
- Functional fault modeling

Motivations

Current situation:

- The efficiency of test generation (quality, speed) is highly depending on
 - the description method (level, language), and
 - fault models
- Because of the growing complexity of systems, gate level methods have become obsolete
- High-Level methods for diagnostic modeling are today emerging, however they are not still mature

Main disadvantages:

- The known methods for fault modeling are
 - dedicated to special classes (i.e. for microprocessors, for RTL, VHDL etc. languages...), not general
 - not well defined and formalized

Contributions

- **High-Level Decision Diagrams** are proposed for diagnostic modeling of digital systems
- A novel DD-based **node fault model** is proposed
- The fault model is **simple** and **formalized**
- Traditional high-level fault models for different abstraction levels of digital systems can be replaced by the new **uniform** fault model
- As the result,
 - the complexity of fault representation is reduced, and
 - the speed of test generation and fault simulation can be increased

Register Level Fault Models

RTL statement:

K: (If T,C) $R_D \leftarrow F(R_{S1}, R_{S2}, \dots R_{Sm}), \rightarrow N$

Components (variables)
of the statement:

K - label
T - timing condition
C - logical condition
 R_D - destination register
 R_S - source register
F - operation (microoperation)
 \leftarrow - data transfer
 $\rightarrow N$ - jump to the next statement

RT level faults:

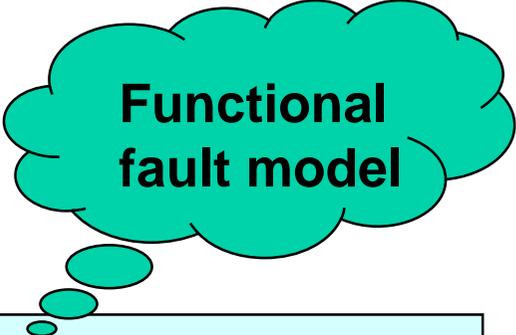
$K \rightarrow K'$ - label faults
 $T \rightarrow T'$ - timing faults
 $C \rightarrow C'$ - logical condition faults
 $R_D \rightarrow R_D$ - register decoding faults
 $R_S \rightarrow R_S$ - data storage faults
 $F \rightarrow F'$ - operation decoding faults
 \leftarrow - data transfer faults
 $\rightarrow N$ - control faults
 $(F) \rightarrow (F)'$ - data manipulation faults

Fault Models and Tests

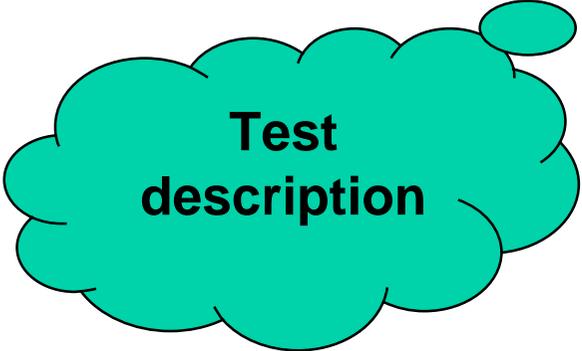
Dedicated functional fault model for multiplexer:

- stuck-at-0 (1) on inputs,
- another input (instead of, additional)
- value, followed by its complement

- value, followed by its complement on a line whose address differs in one bit

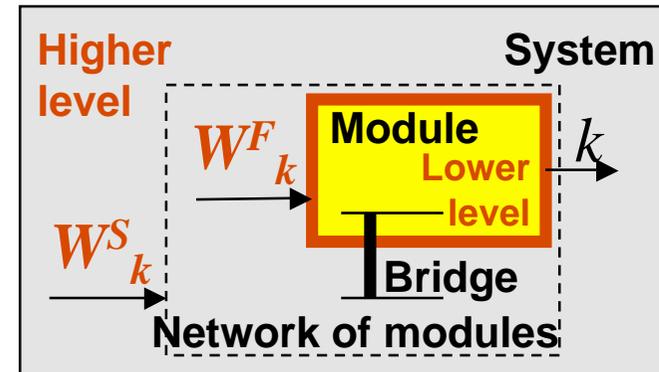
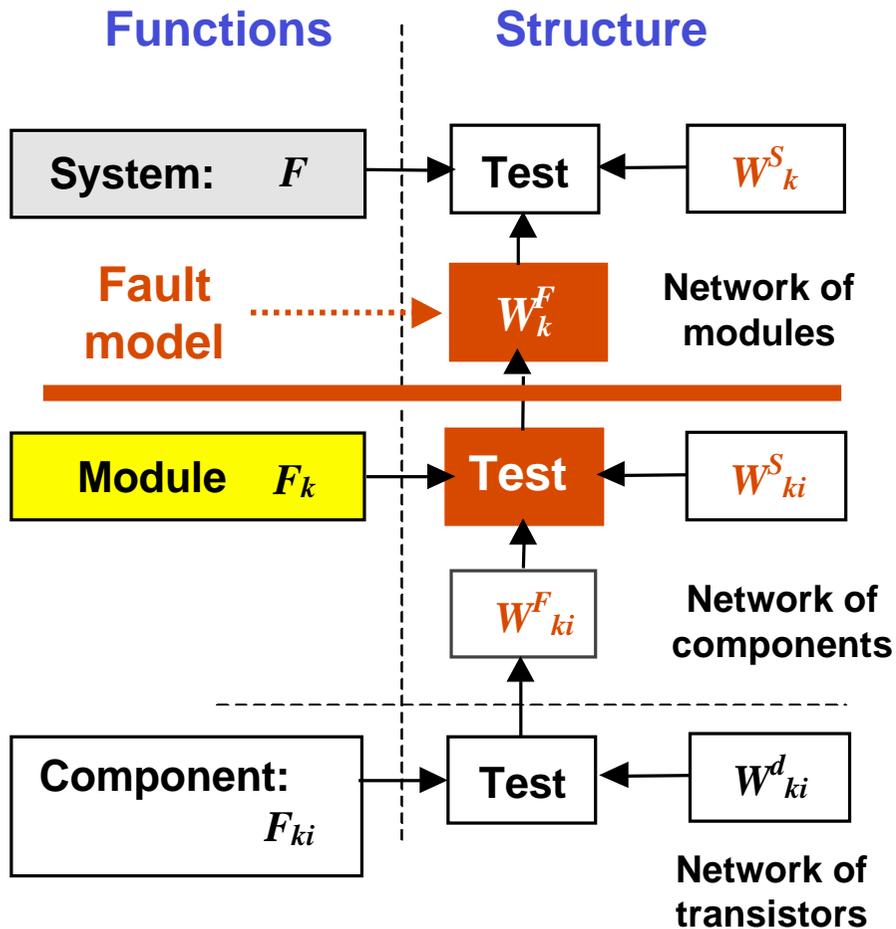


Functional
fault model



Test
description

Hierarchical Fault Modeling



W_k^F interpretation:

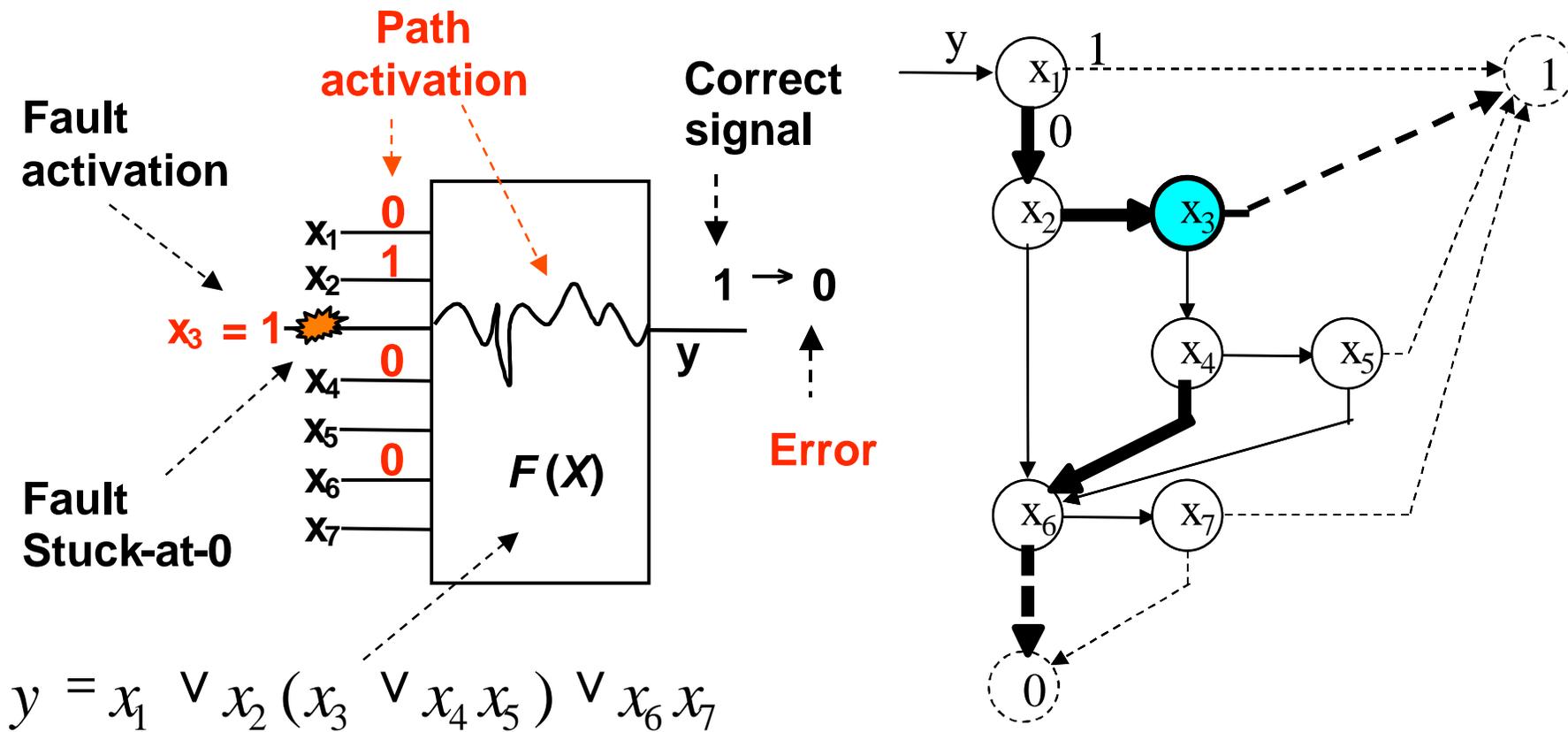
Test – at the lower level

Fault model – at the higher level

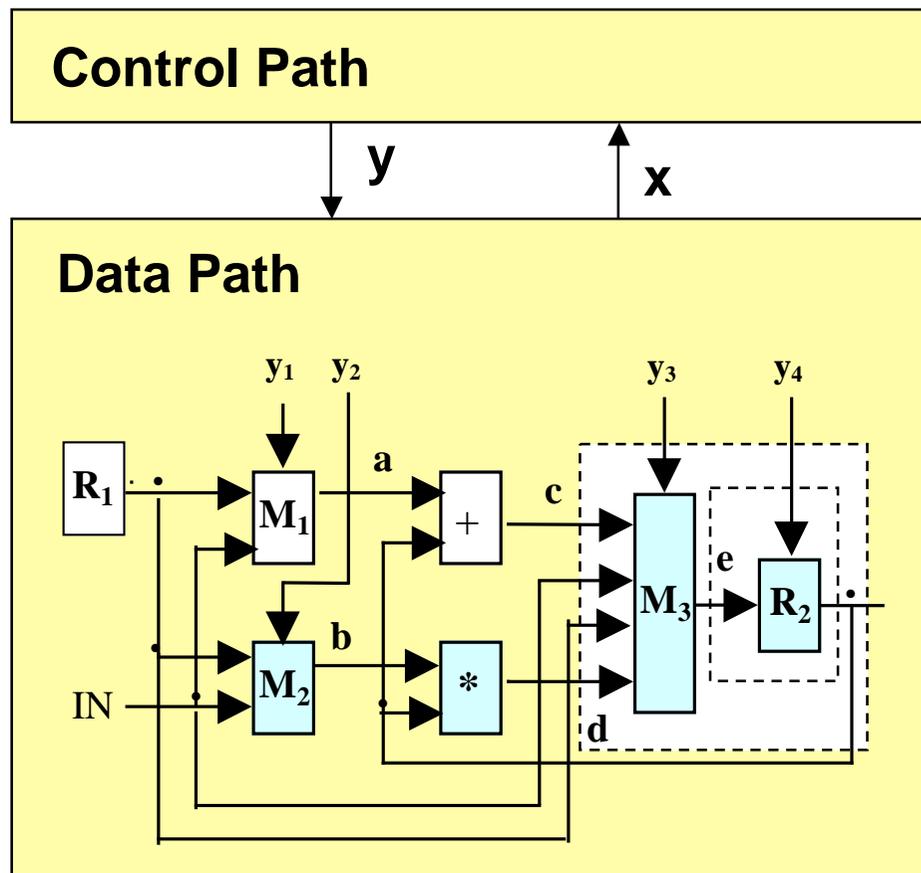
Interface between levels

Logic Level Faults on SSBDDs

Fault modeling on Structurally Synthesized BDDs:



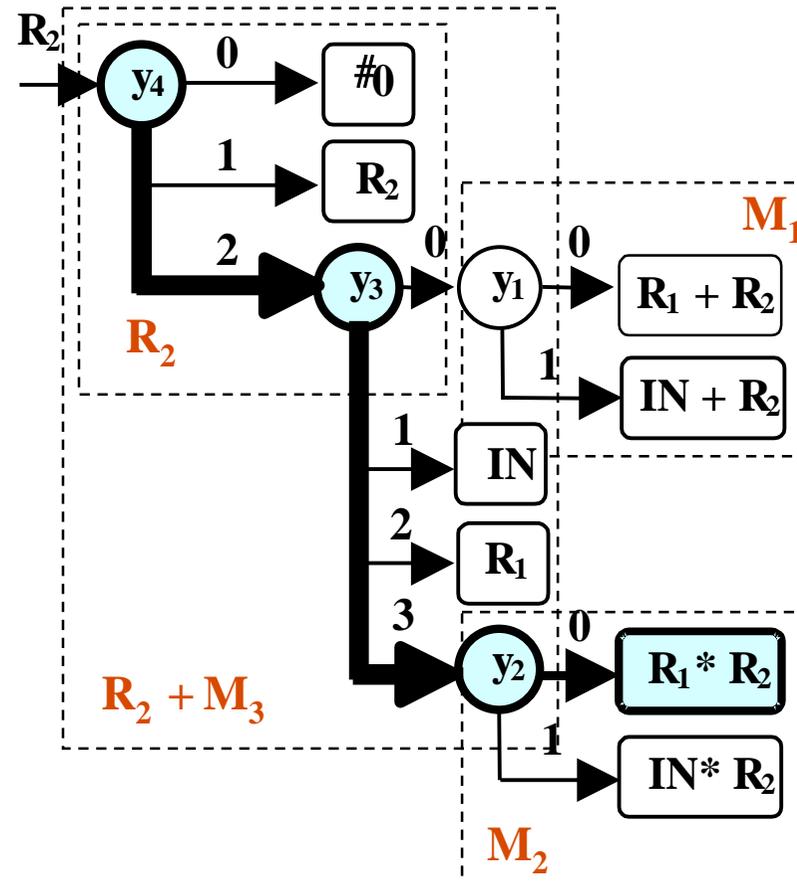
Data Path in Digital Systems



M_1		
y_1	Function	
0	$M_1 = R_1$	
1	$M_1 = IN$	
M_2		
y_2	Function	
0	$M_2 = R_1$	
1	$M_2 = IN$	
M_3		
y_3	Function	
0	$M_3 = M_1 + R_2$	
1	$M_3 = IN$	
2	$M_3 = R_1$	
3	$M_3 = M_2 * R_2$	
R_2		
y_4	Operation	Function
0	Reset	$R_2 = 0$
1	Hold	$R_2 = R_2$
2	Load	$R_2 = M_3$

Decision Diagram of the Data Path

M_1		
y_1	Function	
0	$M_1 = R_1$	
1	$M_1 = IN$	
M_2		
y_2	Function	
0	$M_2 = R_1$	
1	$M_2 = IN$	
M_3		
y_3	Function	
0	$M_3 = M_1 + R_2$	
1	$M_3 = IN$	
2	$M_3 = R_1$	
3	$M_3 = M_2 * R_2$	
R_2		
y_4	Operation	Function
0	Reset	$R_2 = 0$
1	Hold	$R_2 = R_2$
2	Load	$R_2 = M_3$



Faults and High-Level Decision Diagrams

RTL-statement:

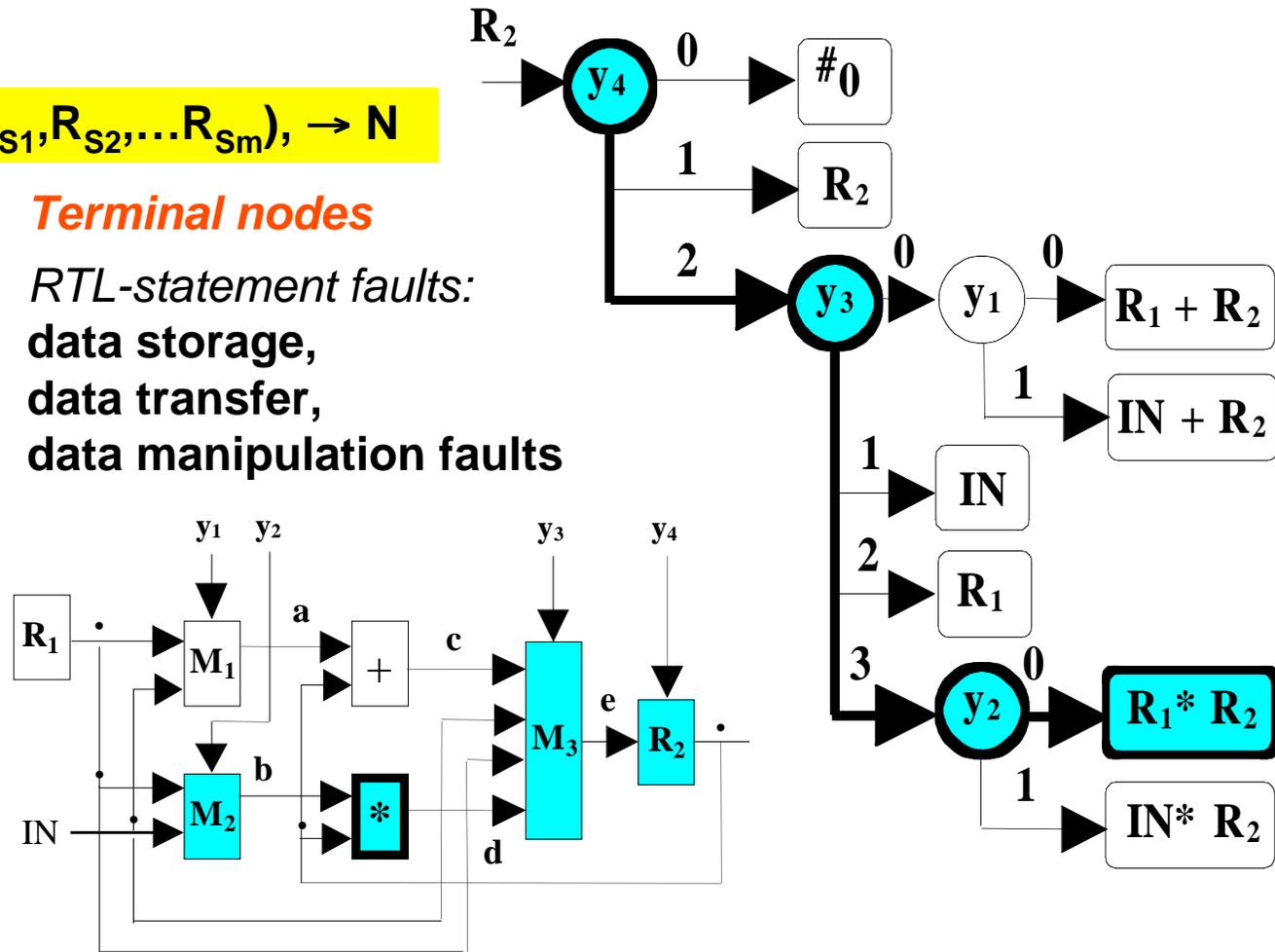
K: (If T,C) $R_D \leftarrow F(R_{S1}, R_{S2}, \dots, R_{Sm}), \rightarrow N$

Nonterminal nodes

RTL-statement faults:
label,
timing condition,
logical condition,
register decoding,
operation decoding,
control faults

Terminal nodes

RTL-statement faults:
data storage,
data transfer,
data manipulation faults



Faults and High-Level Decision Diagrams

RTL-statement:

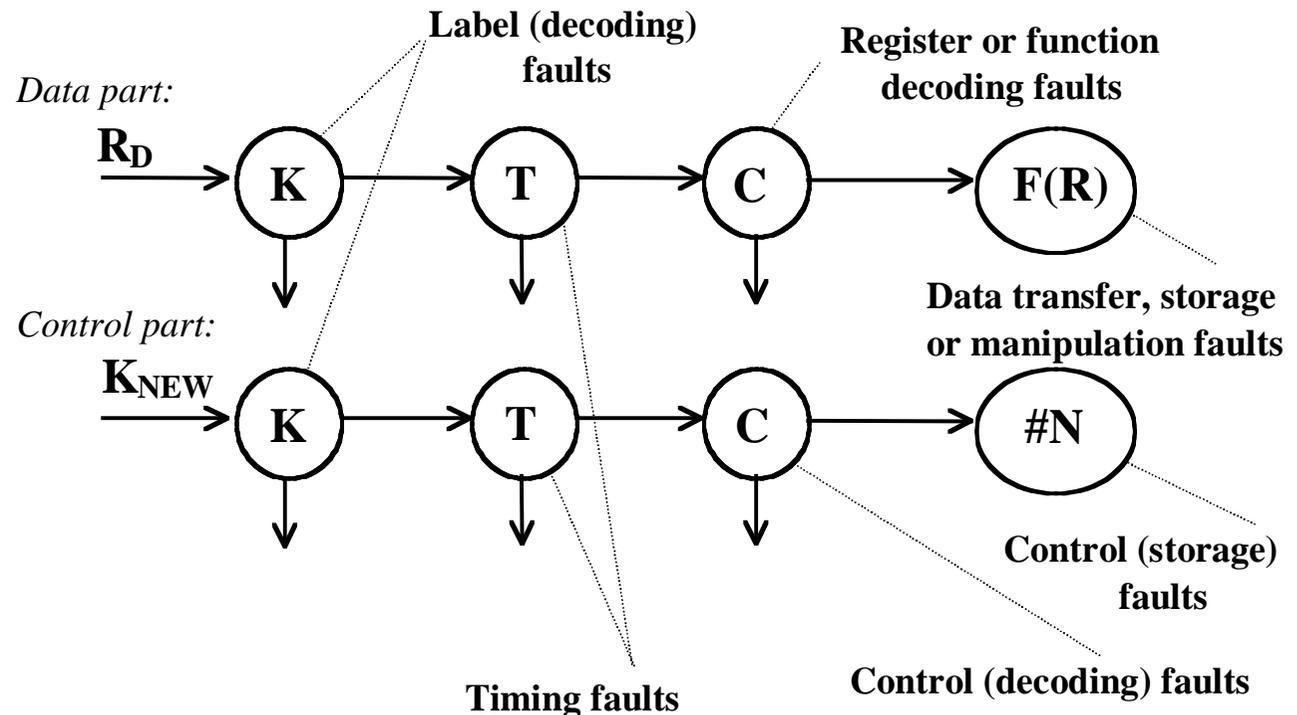
$$K: (If T, C) R_D \leftarrow F(R_{S1}, R_{S2}, \dots, R_{Sm}), \rightarrow N$$

Nonterminal nodes

RTL-statement faults:
 label,
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Terminal nodes

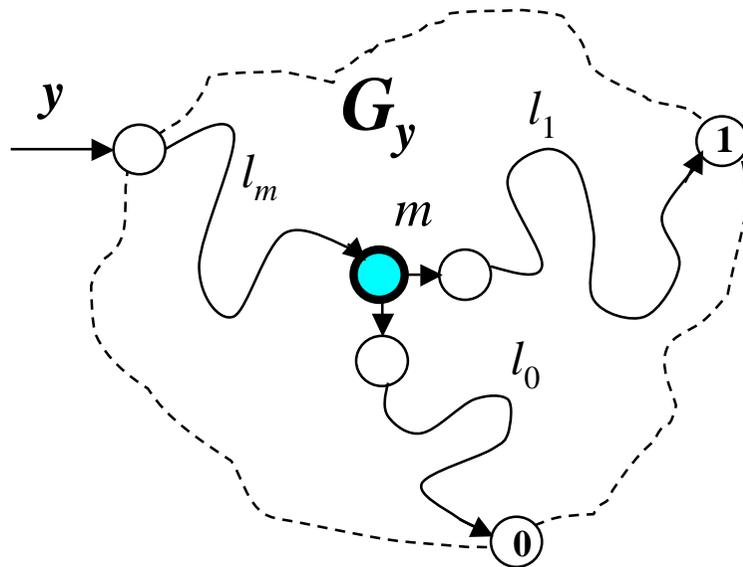
RTL-statement faults:
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Fault Modeling on DDS

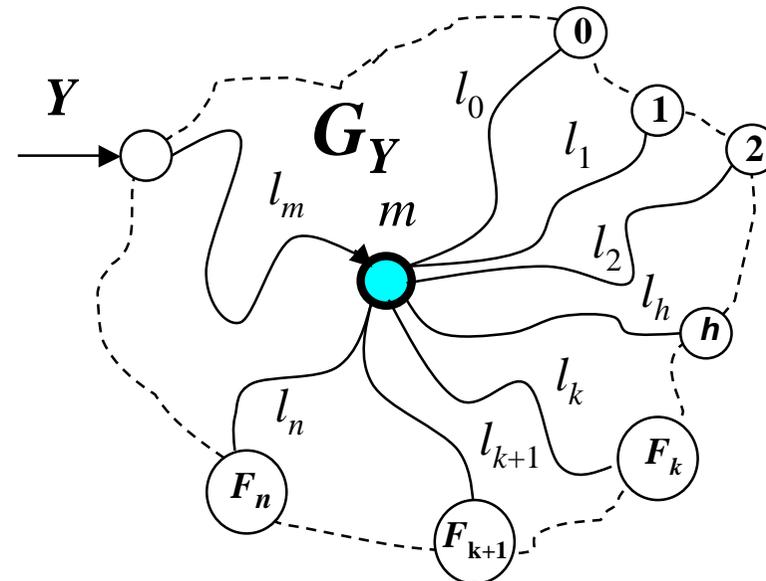
Binary DD

with 2 terminal nodes and
2 outputs
from each node



General case of DD

with $n \geq 2$ terminal nodes and
 $n \geq 2$ outputs
from each node



Fault Model for Decision Diagrams

- Each path in a DD describes the behavior of the system in a specific **mode** of operation
- **The faults** having effect on the behaviour can be associated with nodes **along the path**
- A fault causes **incorrect leaving** the path activated by a test

Fault Model for Decision Diagrams

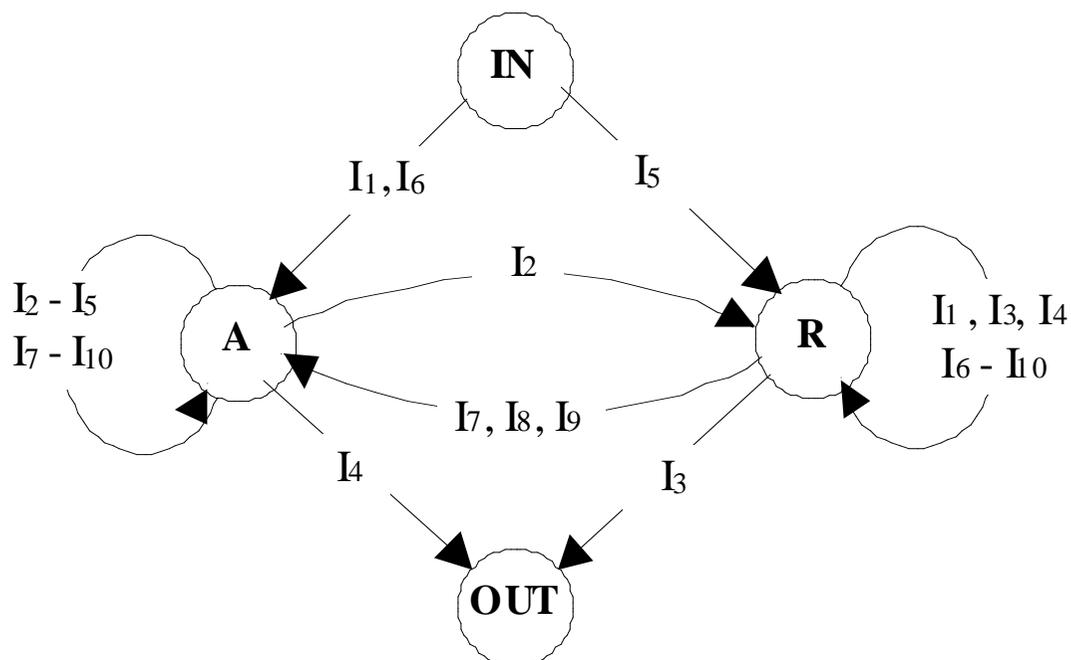
- D1:** the output edge for $x(m) = i$ of a node m is **always activated**
- D2:** the output edge for $x(m) = i$ of a node m is **broken**
- D3:** instead of the given edge, **another edge** or a set of edges is **activated**

Microprocessor Modeling with S-Graphs

Instruction set of a Microprocessor:

I ₁ :	MVI A,D	$A \leftarrow IN$
I ₂ :	MOV R,A	$R \leftarrow A$
I ₃ :	MOV M,R	$OUT \leftarrow R$
I ₄ :	MOV M,A	$OUT \leftarrow A$
I ₅ :	MOV R,M	$R \leftarrow IN$
I ₆ :	MOV A,M	$A \leftarrow IN$
I ₇ :	ADD R	$A \leftarrow A + R$
I ₈ :	ORA R	$A \leftarrow A \vee R$
I ₉ :	ANA R	$A \leftarrow A \wedge R$
I ₁₀ :	CMA A,D	$A \leftarrow \neg A$

S-Graph:



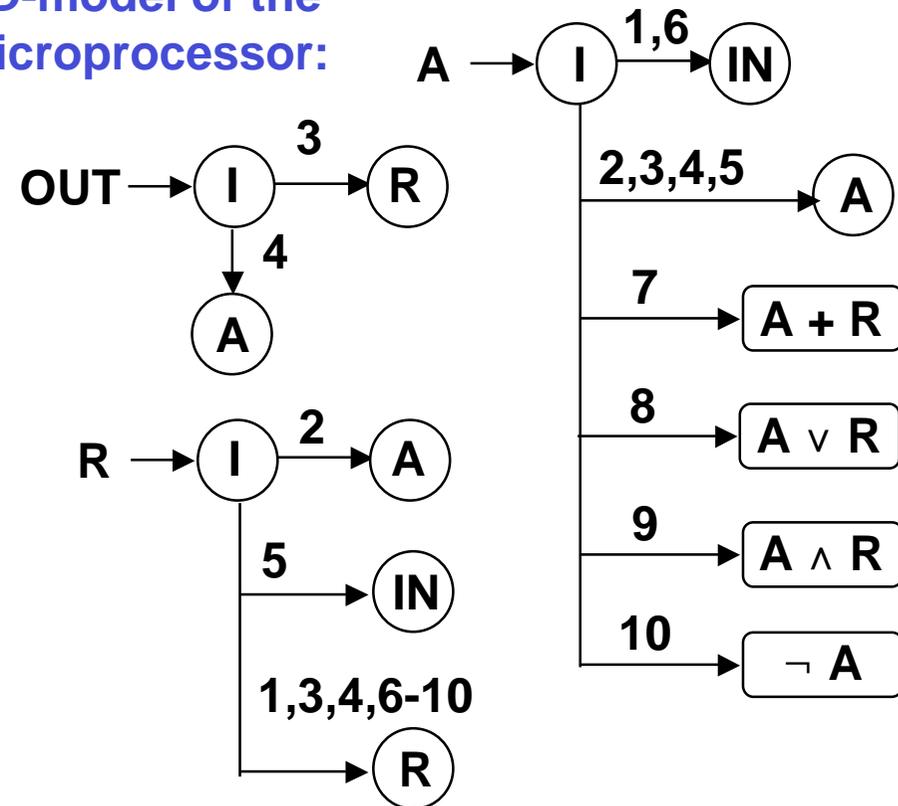
Test Generation for Microprocessors

High-Level DDs for a microprocessor (example):

Instruction set:

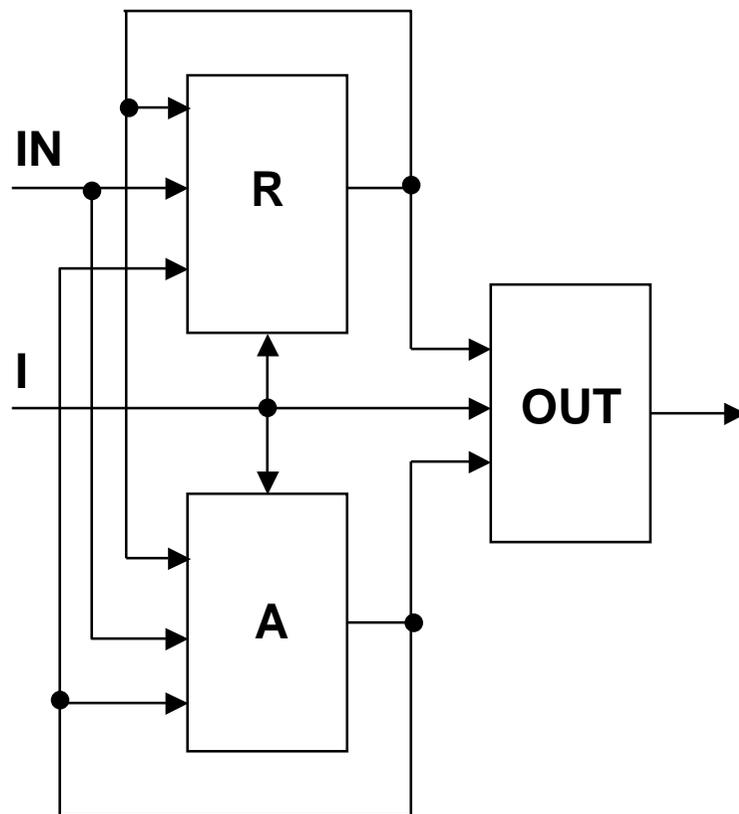
I ₁ :	MVI A,D	A ← IN
I ₂ :	MOV R,A	R ← A
I ₃ :	MOV M,R	OUT ← R
I ₄ :	MOV M,A	OUT ← A
I ₅ :	MOV R,M	R ← IN
I ₆ :	MOV A,M	A ← IN
I ₇ :	ADD R	A ← A + R
I ₈ :	ORA R	A ← A ∨ R
I ₉ :	ANA R	A ← A ∧ R
I ₁₀ :	CMA A,D	A ← ¬ A

DD-model of the microprocessor:

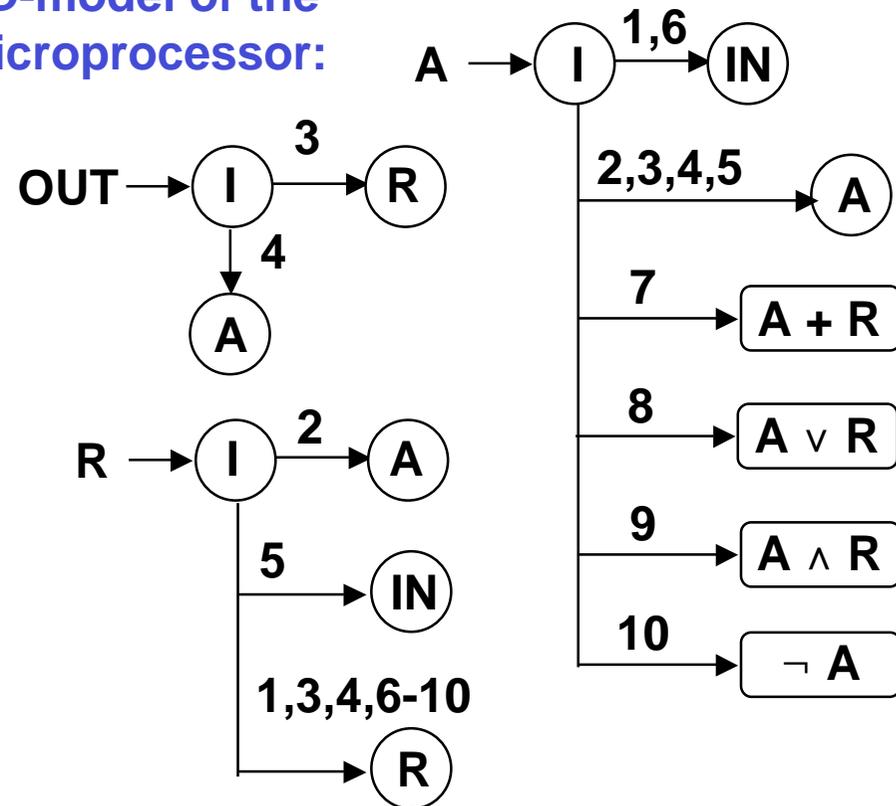


Decision Diagrams for Microprocessors

High-Level DD-based structure of the microprocessor (example):



DD-model of the microprocessor:



Microprocessor Fault Model

Faults affecting the operation of microprocessor can be divided into the following classes:

- **addressing faults affecting register decoding;**
- **addressing faults affecting the instruction decoding and -sequencing functions;**
- **faults in the data-storage function;**
- **faults in the data-transfer function;**
- **faults in the data-manipulation function.**

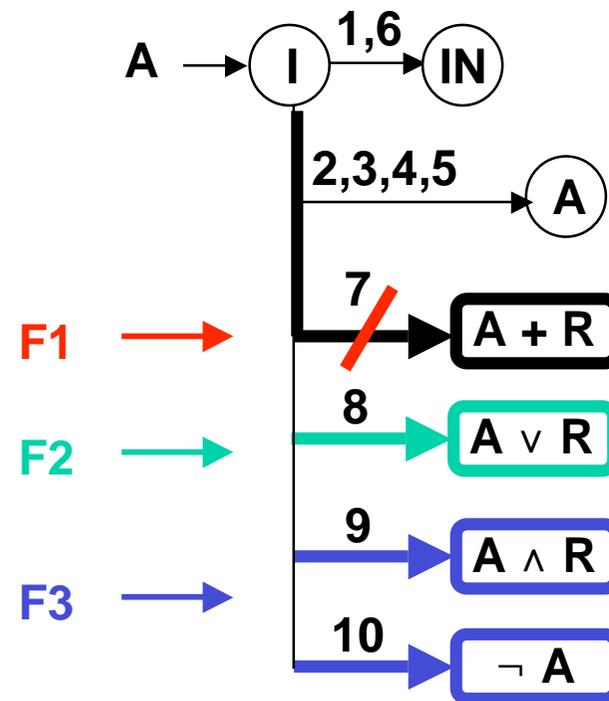
Microprocessor Fault Model

For **multiplexers** under a fault, for a given source address any of the following may happen:

F1: no source is selected

F2: wrong source is selected;

F3: more than one source is selected and the multiplexer output is either a wired-AND or a wired-OR function of the sources, depending on the technology.

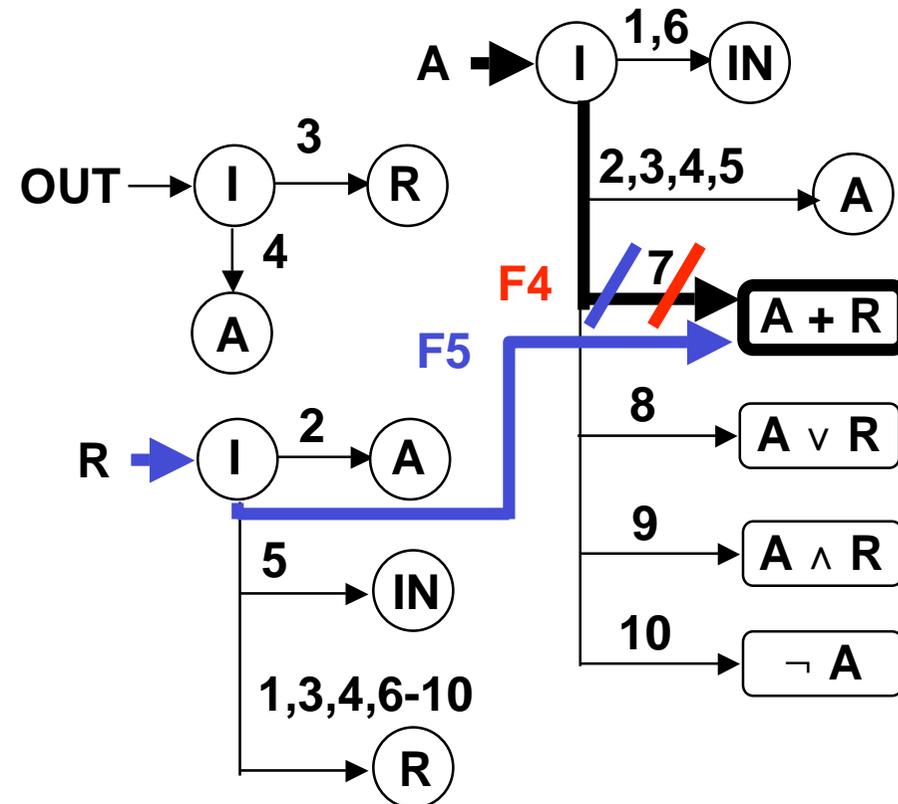


Microprocessor Fault Model

For **demultiplexers** under a fault, for a given destination address:

F4: no destination is selected

F5: instead of, or in addition to the selected correct destination, one or more other destinations are selected



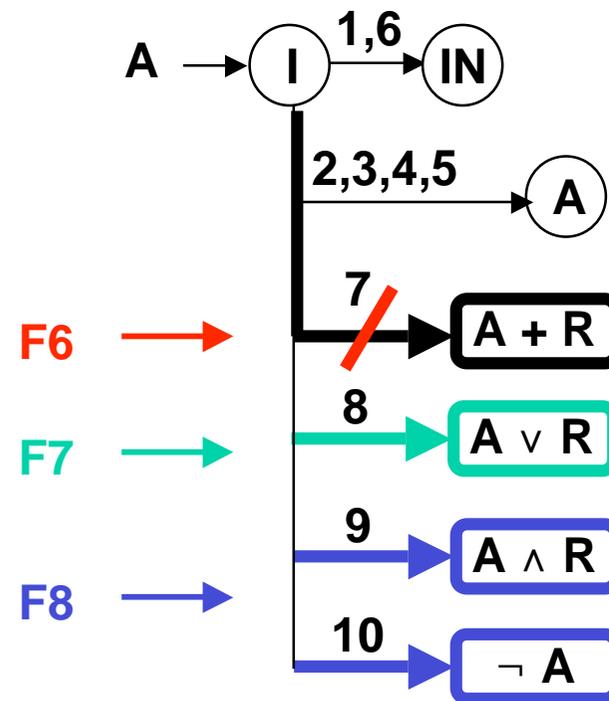
Microprocessor Fault Model

Addressing faults affecting the execution of an instruction may cause the following fault effects:

F6: one or more microorders not activated by the microinstructions of I

F7: microorders are erroneously activated by the microinstructions of I

F8: a different set of microinstructions is activated instead of, or in addition to, the microinstructions of I



Microprocessor Fault Model

The data storage faults:

F9: one or more cells stuck at 0 or 1;

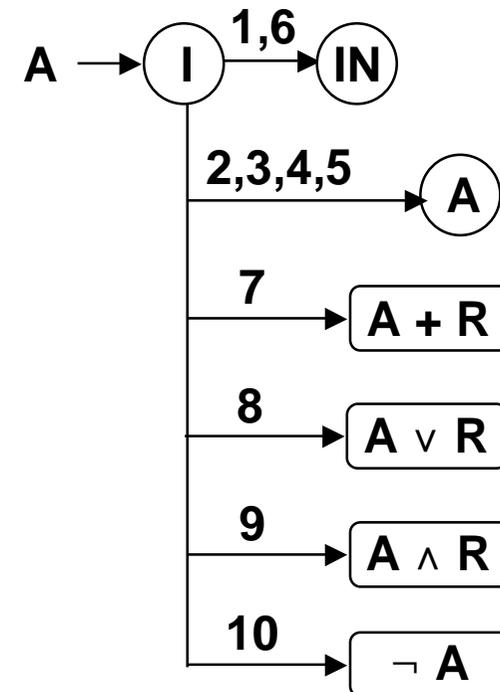
F10: one or more cells fail to make a $0 \rightarrow 1$ or $1 \rightarrow 0$ transitions;

F11: two or more pairs of cells are coupled;

For buses under a fault:

F12: one or more lines stuck at 0 or 1;

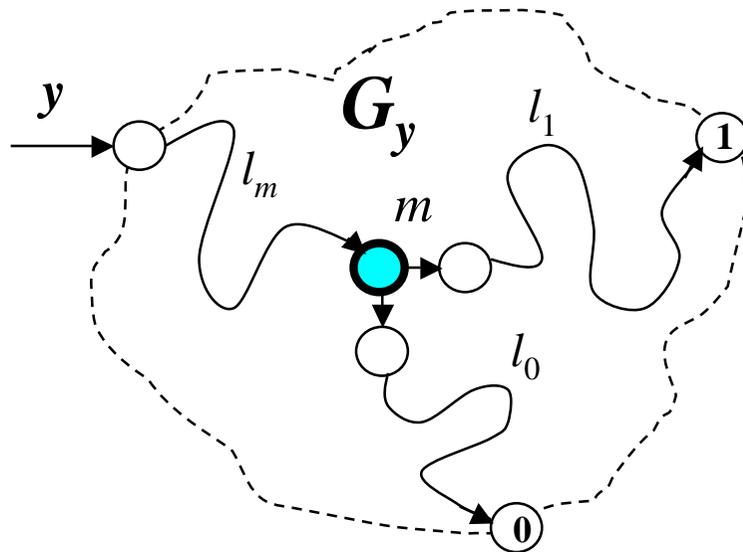
F13: one or more lines form a wired-OR or wired-AND function due to shorts or spurious coupling



Test Generation on DDS

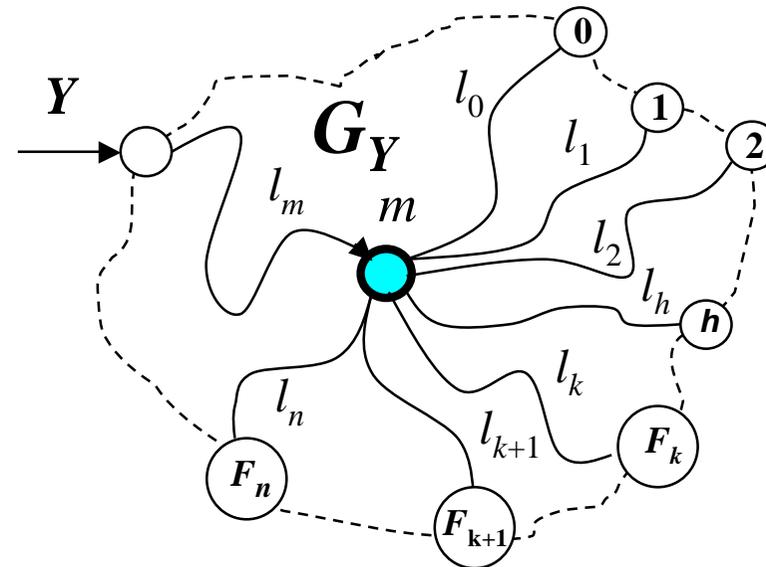
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General case of DD

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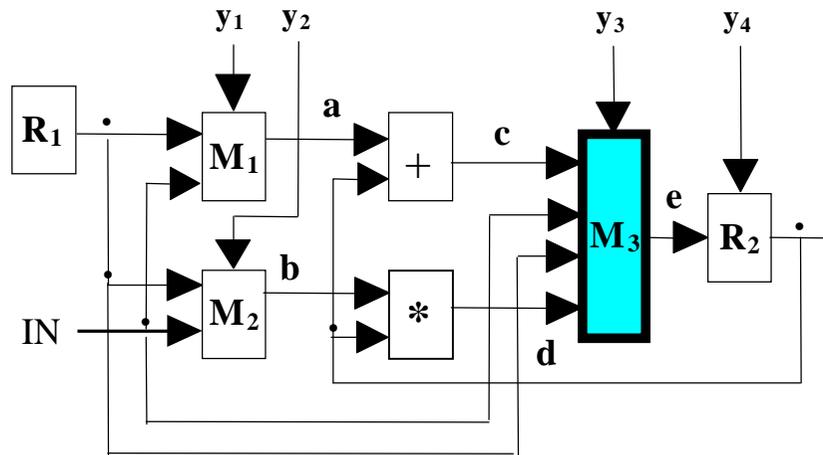


Hierarchical Test Generation on DDs

Hierarchical test generation with DDs: Scanning test

Single path activation in a single DD
Data function $R_1 * R_2$ is tested

Data path

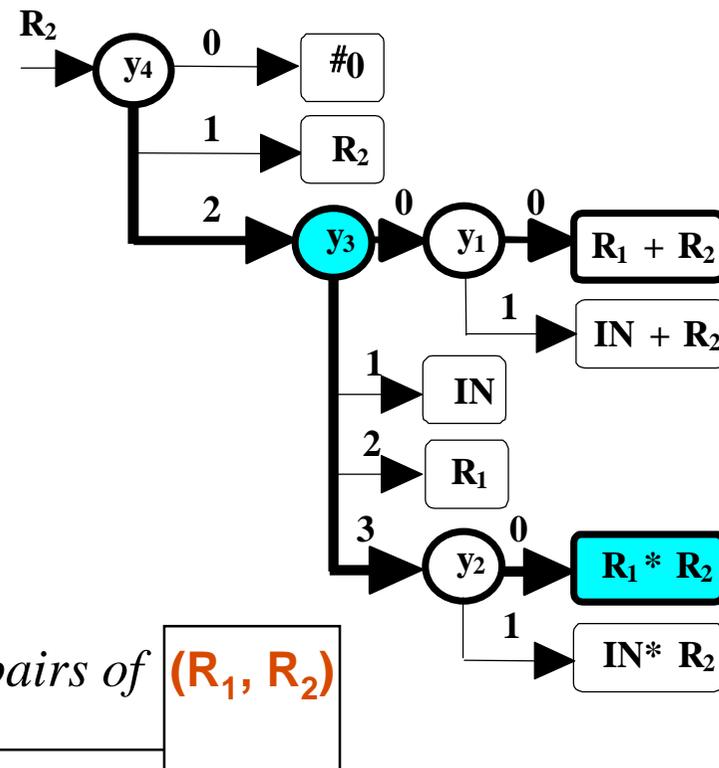


Test program: Control: $y_1 y_2 y_3 y_4 = x032$

Data: For all specified pairs of (R_1, R_2)

Low level test data (constraints W)

Decision Diagram



Test Generation on High Level DDs

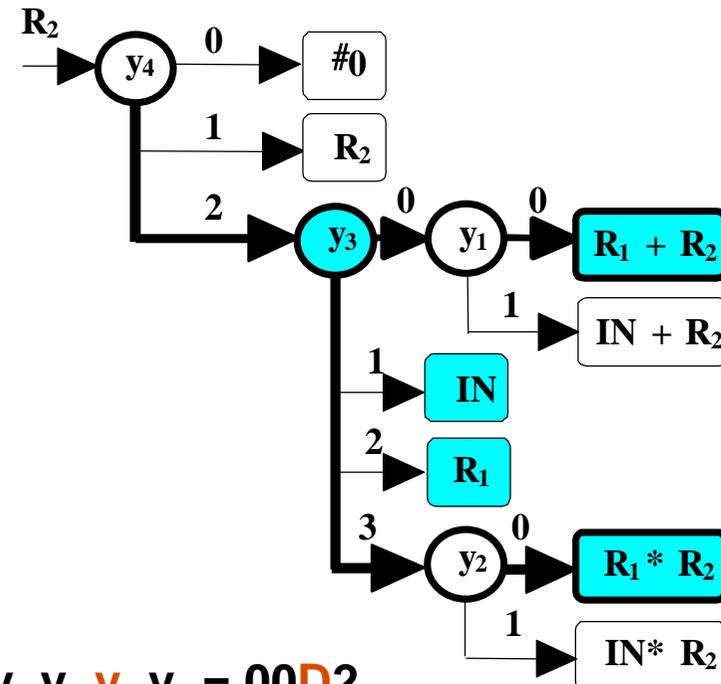
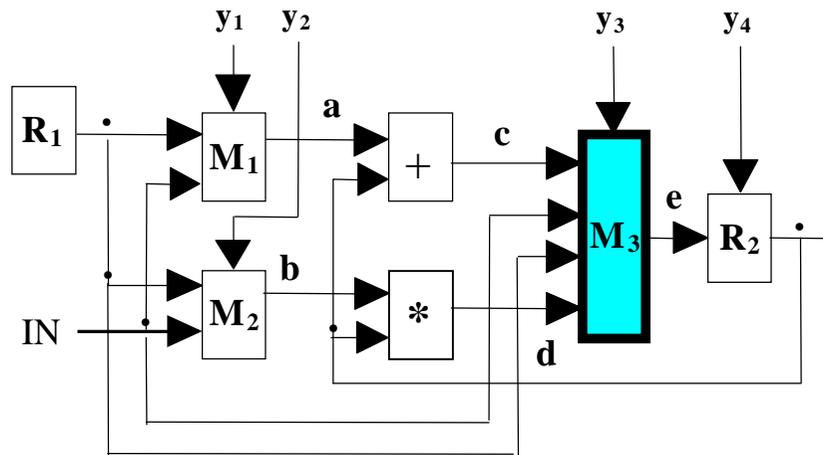
High-level test generation with DDs:

Conformity test

Multiple paths activation in a single DD
Control function y_3 is tested

Decision Diagram

Data path

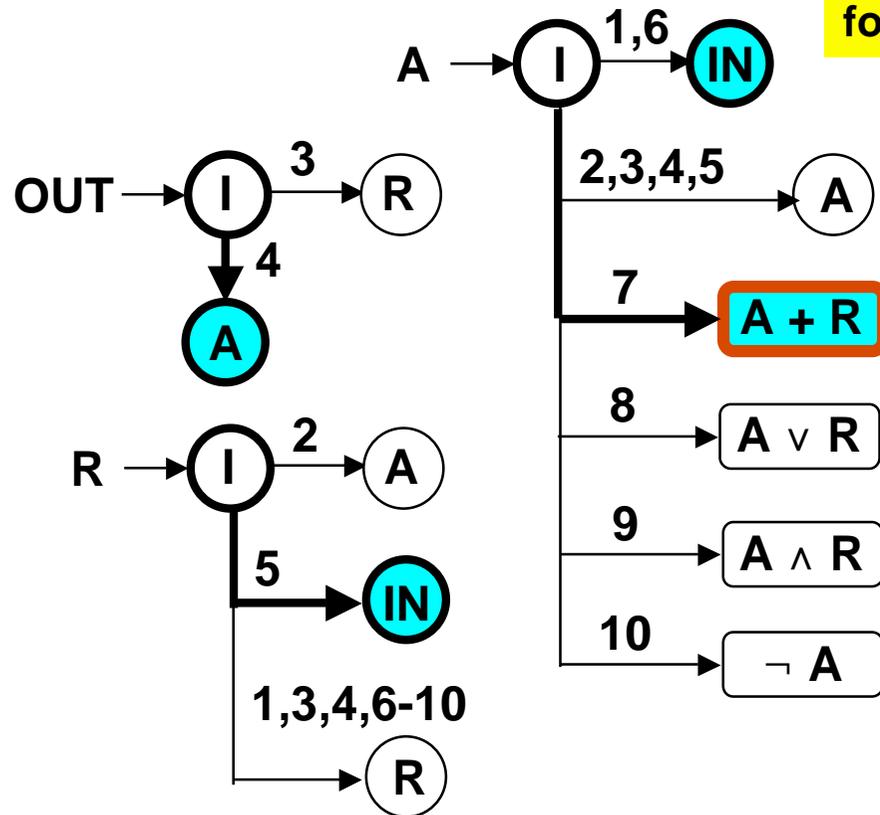


Test program: Control: For $D = 0,1,2,3$: $y_1 y_2 y_3 y_4 = 00D2$

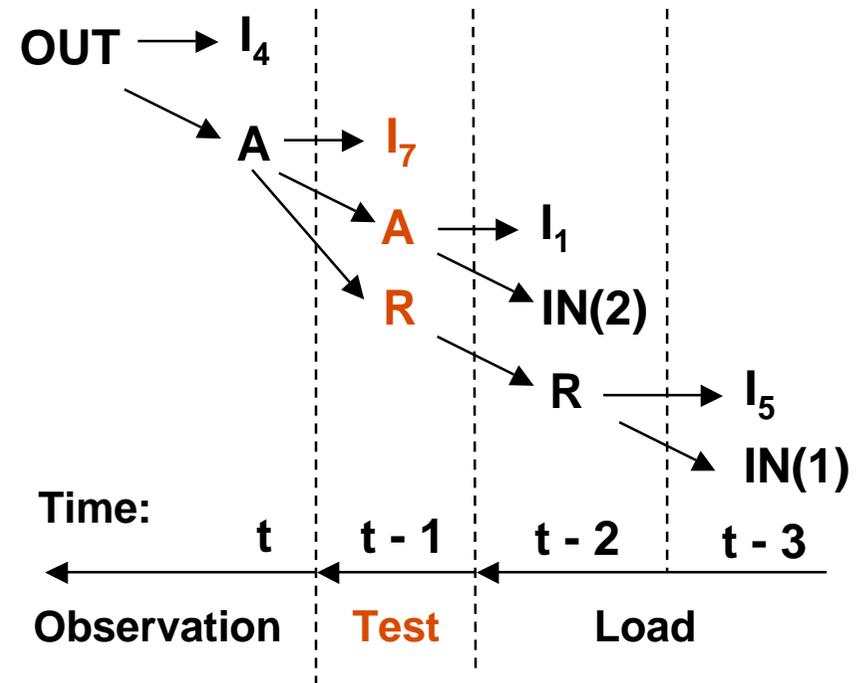
Activating high-level faults: Data: Solution of $R_1 + R_2 \neq IN \neq R_1 \neq R_1 * R_2$

Test Generation for Microprocessors

DD-model of the microprocessor:

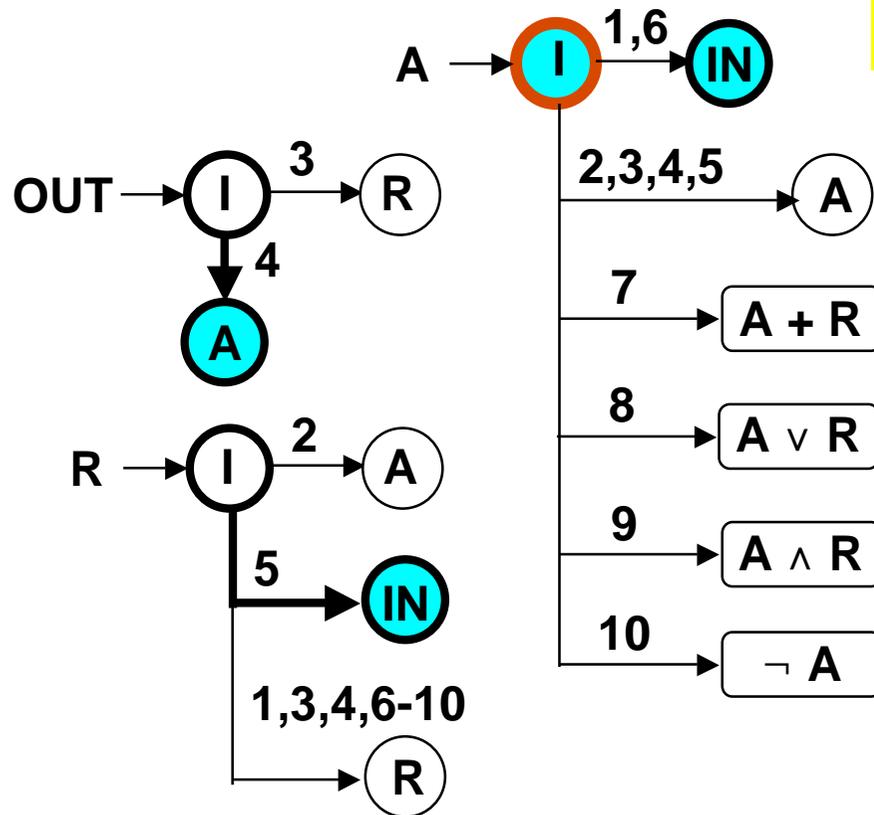


Scanning test program for adder:
 Instruction sequence $T = I_5 (R)I_1 (A)I_7 I_4$
 for all needed pairs of (A,R)

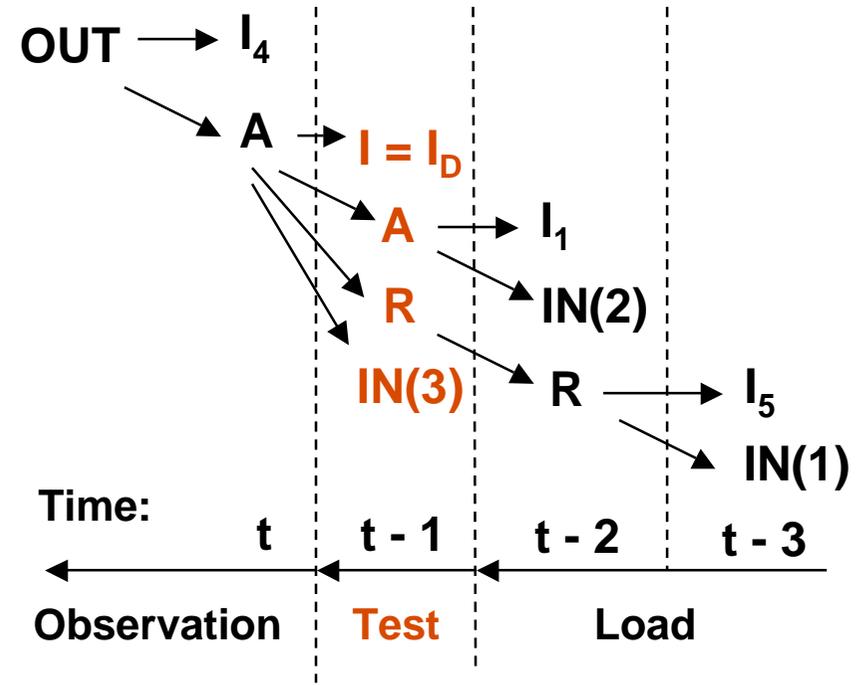


Test Generation for Microprocessors

DD-model of the microprocessor:



Conformity test program for decoding I:
 Instruction sequence $T = I_5 I_1 D I_4$
 for all $D \in \{I_1 - I_{10}\}$ at given A,R,IN(3)



Experimental results

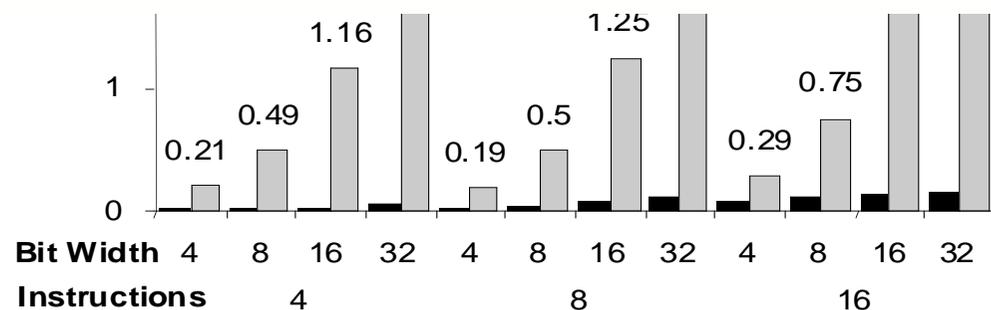
Experimental results with RISC processors

BW	ATPG	Time	Test	Faults
4	HTPG	0,08	224	900
	Synopsys	0,29	46	855
8	HTPG	0,10	224	1612
	Synopsys	0,75	64	1531
16	HTPG	0,13	224	3016
	Synopsys	1,86	73	2861
32	HTPG	0,15	224	5908
	Synopsys	5,57	84	5607

HTPG – high level

Synopsys – gate level

Gate-level fault coverage – 100%



Conclusions

- Different fault models for different representation levels of digital systems can be replaced on **DDs** by the **uniform node fault model**
- It allows to represent groups of structural faults through groups of functional faults
- As the result, the **complexity** of fault representation can be reduced, and the **simulation speed** can be raised
- The fault model on DDs can be regarded as a **generalization**
 - of the classical gate-level stuck-at fault model, and
 - of the known higher level fault models

www.pld.ttu.ee/~raiub/