

RATP safety approach for railway signalling systems

ReSIST summer School 2007

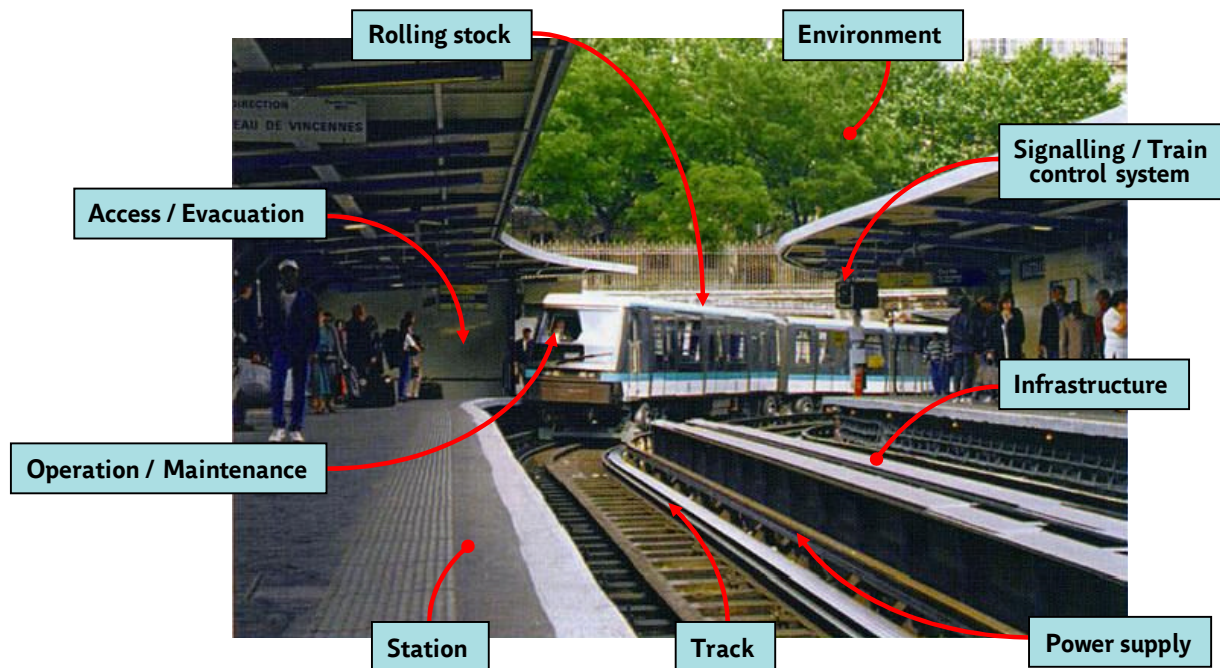
Pierre CHARTIER



Summary

1. Introduction
2. Hardware fault detection
3. Software fault avoidance

Global railway system

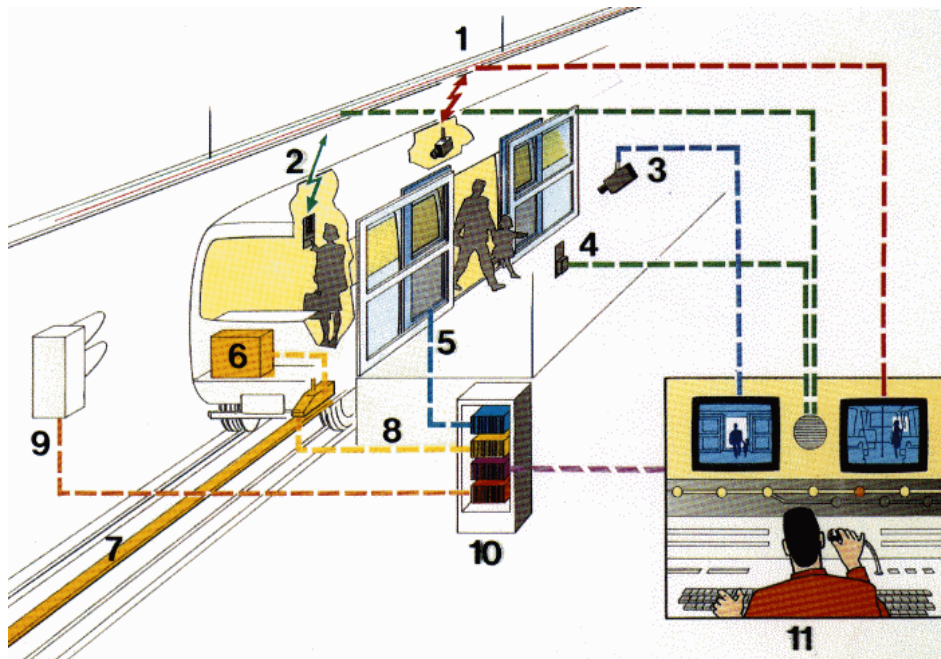


Events to be feared at global system level

- Fire / explosion
- Derailment / overturning
- Panic
- Electrocution / burn
- Collision
- Individual accidents (fall, ...)
- Others (terrorist attack, natural disaster, structure breaking, ...)



Transport system overview (METEOR example)



Railway signalling system

Main protection against collision and derailment

- Safety critical mission

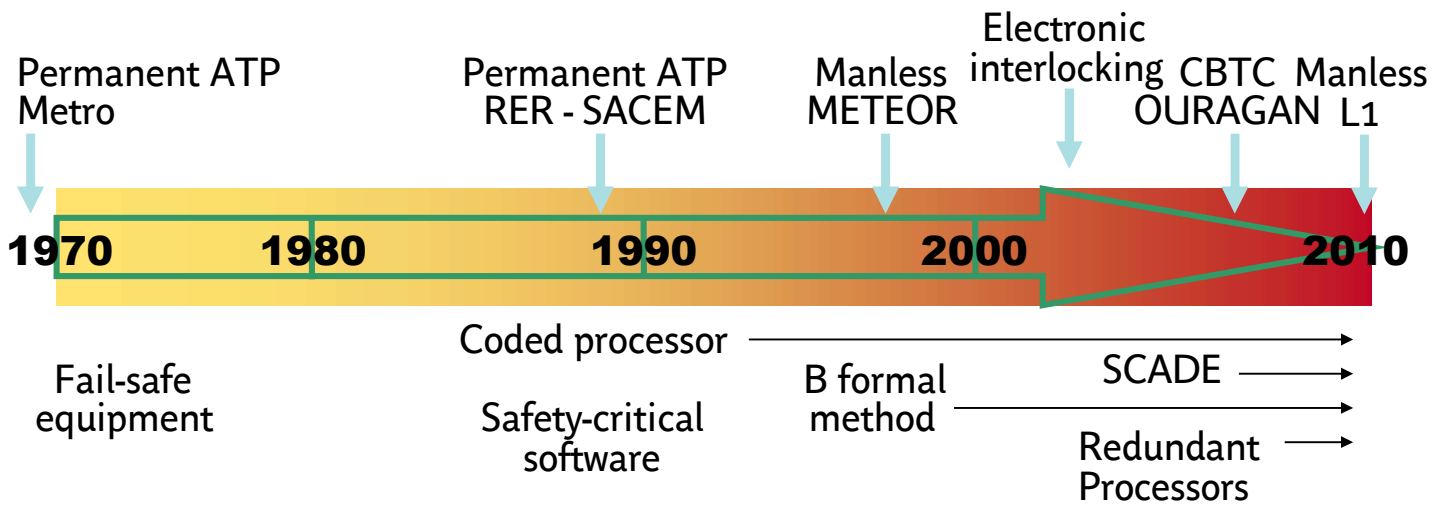
Historically 2 types of system

- Interlocking
- Automatic train protection

Main safety measure : stop all trains and power off the traction power supply



RATP technical evolution



SACEM – RER A

Saturation of RER line A (80's) → SACEM project

- Objective :
To increase transport offer by raising train frequency
- But incompatibility between
 - train spacing reduction,
 - and traditional signaling



SACEM – RER A

Automatic Train Protection

- Control train spacing
- Control train speed
- Protect switching zones
- Switch between cab signal and trackside signalling

→ First safety-critical computing system in railways

SACEM – RER A

SACEM functions require the use of computers

Two main concerns :

- Detection of errors due to hardware
→ coded processor
- Avoiding faults in software
→ formal methods



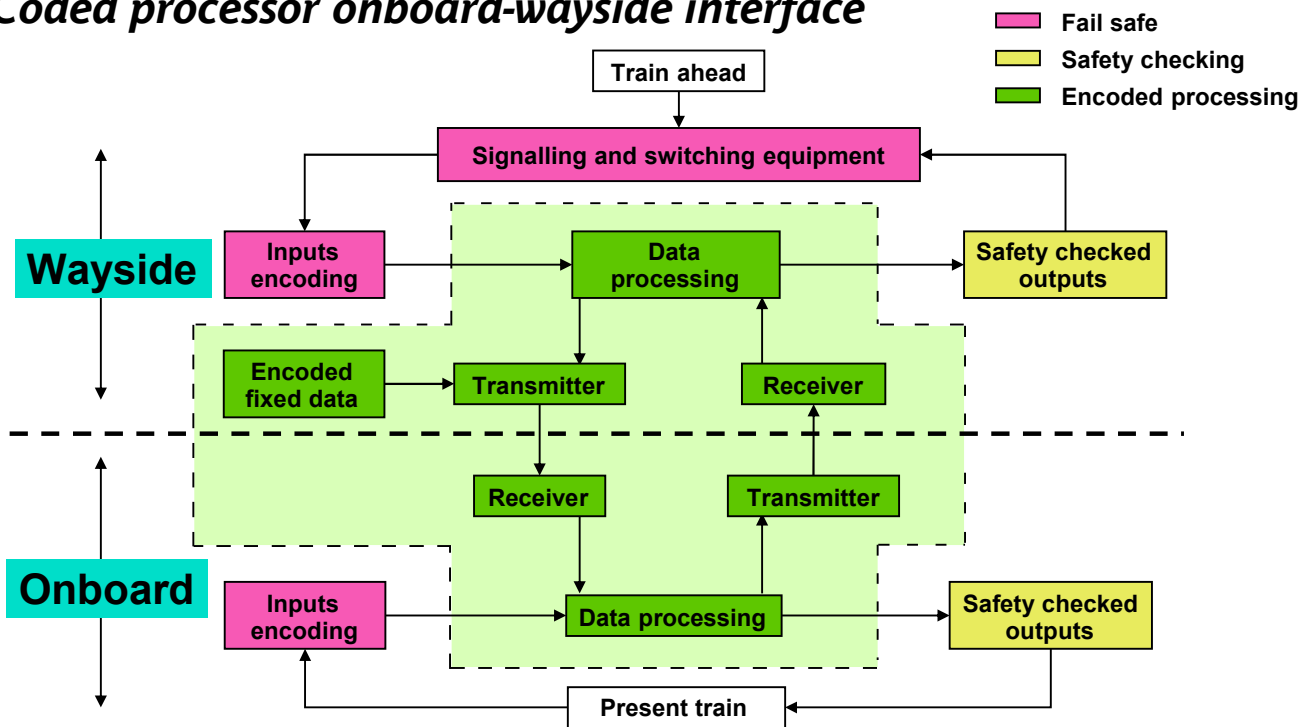
Hardware fault detection

Hardware fault detection

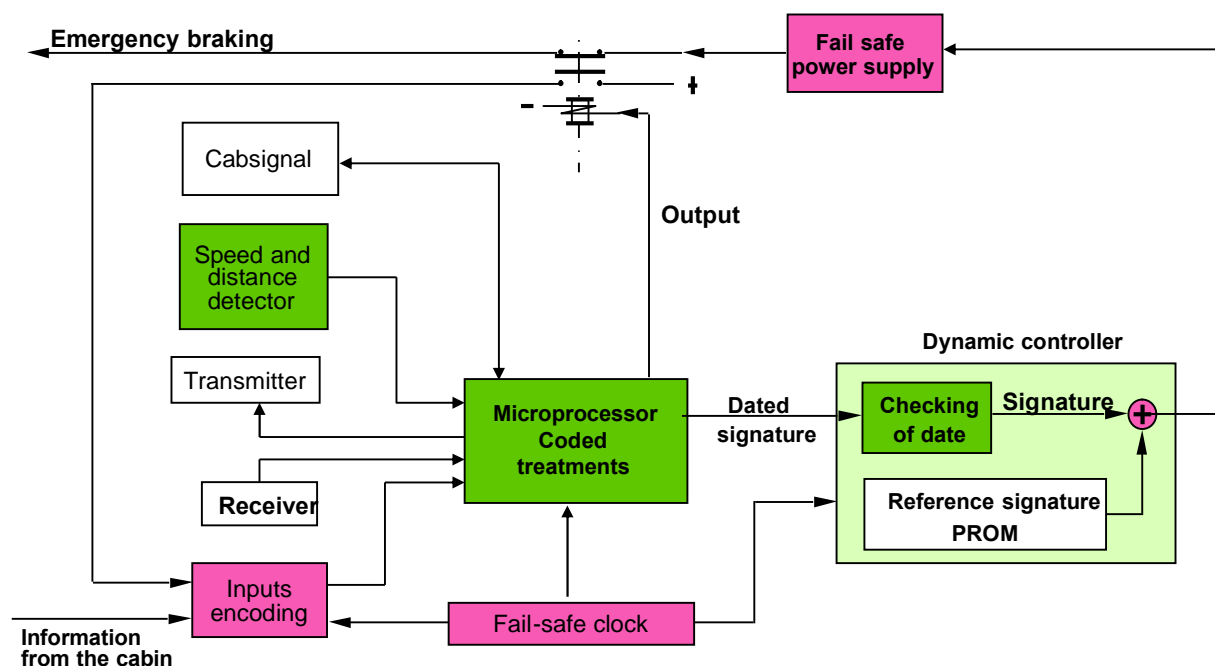
Coded processor – main concepts

- Based on data and program encoding
- Encoding done automatically by specific tools
- Detect run-time errors
- If an error is detected, the hardware sets the system in a fail-safe state

Coded processor onboard-wayside interface



Coded processor



Coded processor – safety data encoding

Data X = functional part X.F and coded part X.C

X.F : N_F bits and X.C : N_C bits

Tasks for the computer

- Acquisition and coding of the fail-safe inputs
- Processing the coded data
- Conversion of coded data into fail-safe outputs
- Setting the system into restrictive state in case of failure



Coded processor – detected errors (1)

Different kinds of errors :

- Arithmetical error
- Operator error
- Operand error
- Memory « non-refreshed » error
- Branch error

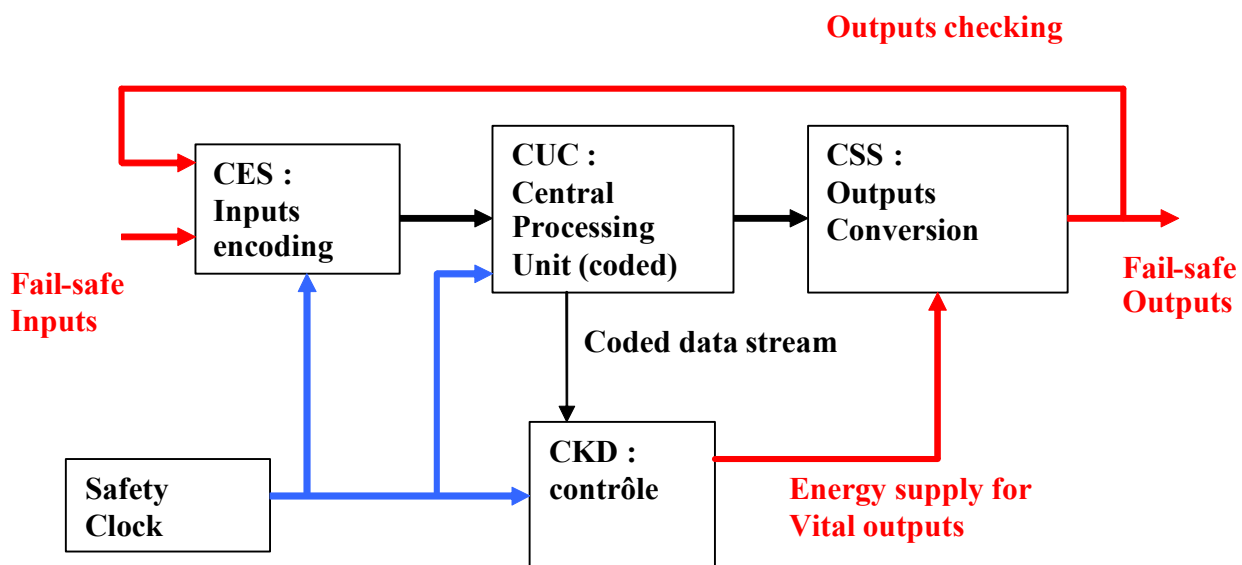


Coded processor – detected errors (2)

Components of the coded part

- Arithmetical error ==> remainder r_{kx}
- Operator error ==> signature B_x
- Operand error ==> signature B_x
- memory « non refreshed » error ==> date D
- Branch error ==> compensation, tracer

Coded processor – architecture



Coded processor – safety integrity level

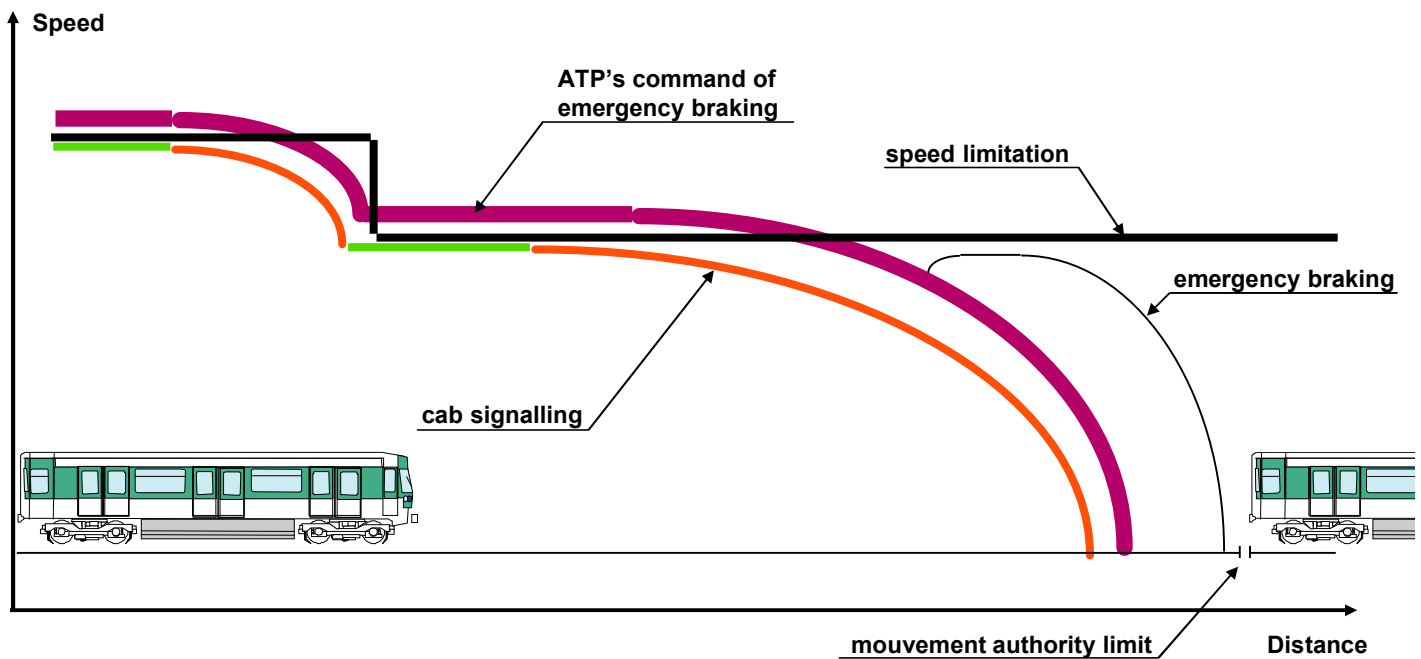
Theoretical result:

- Coded processor alone $\rightarrow 10^{-12} \text{ h}^{-1}$
- Including transmission between onboard and trackside equipment $\rightarrow 10^{-9} \text{ h}^{-1}$

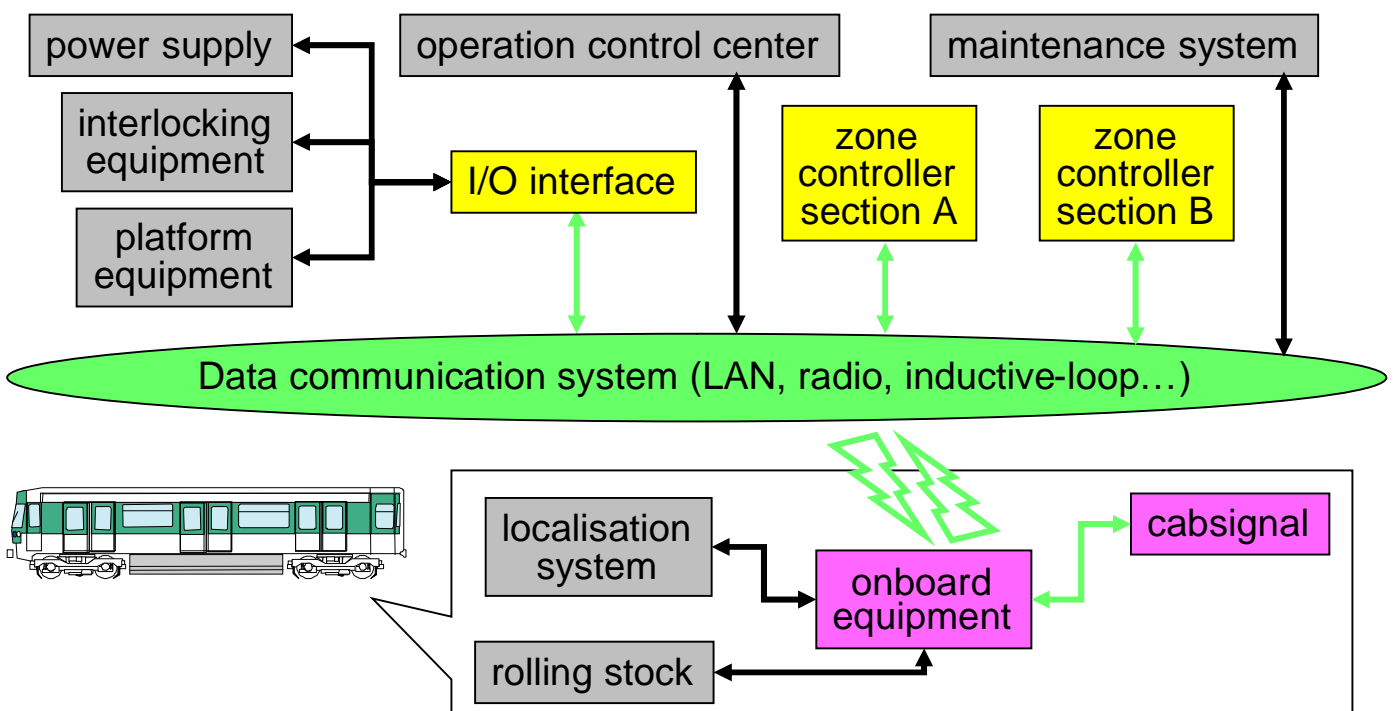


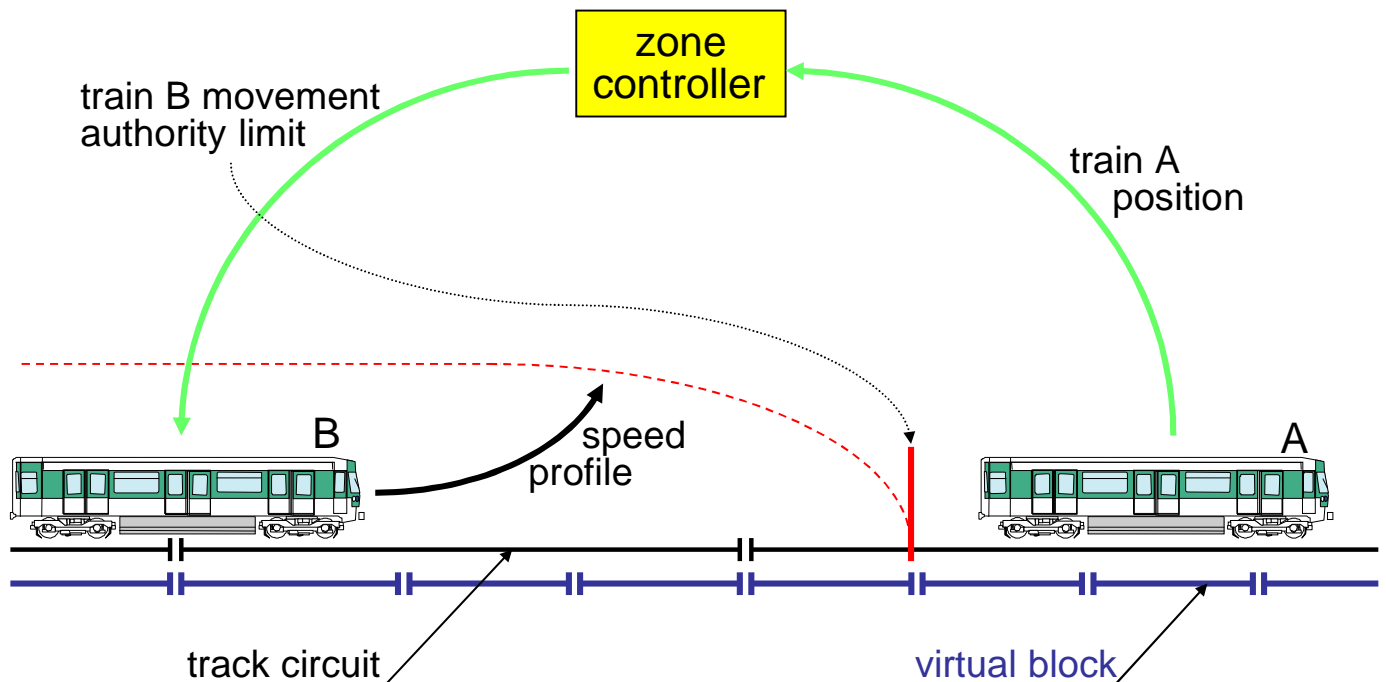
Software fault avoidance

ATP role



Communication-based train control (CBTC) systems



CBTC operation**Communication-based train control (CBTC) systems**

Automatic Train Protection (ATP)

Automatic Train Operation (ATO)

Automatic Train Supervision (ATS)

Formal methods

1988 SACEM - First safety software in railways

- Usual (unformal) software specification issues
 - lack of global approach with the system designer point of view
 - ambiguous, not legible, not coherent, not complete
- Validation issues
 - no certitude that the fonctionnal tests are sufficient

1998 First run of the subway line 14 Météor

The B method is used to obtain :

- a reliable and exact software design from specifications to runtime code



B formal method

Goal

- To get a software which meets completely its fonctionnal specification by construction

Application fields

- Sequential code with no interruptions (real time aspects, low level softwares, operating kernels are not taken into account)

Large spectrum language

- Unified framework and continuous process from specification to code



B formal method

High level language

- Abstract operators for specification needs
- Concrete instructions similar to ADA or C one's

Model oriented approach

- Software = data + properties + operations

Refinement process

- Translation of the abstract machines into concrete modules, and finally into code

Proof obligations

- Conditions to check to ensure a strict mathematical construction of softwares

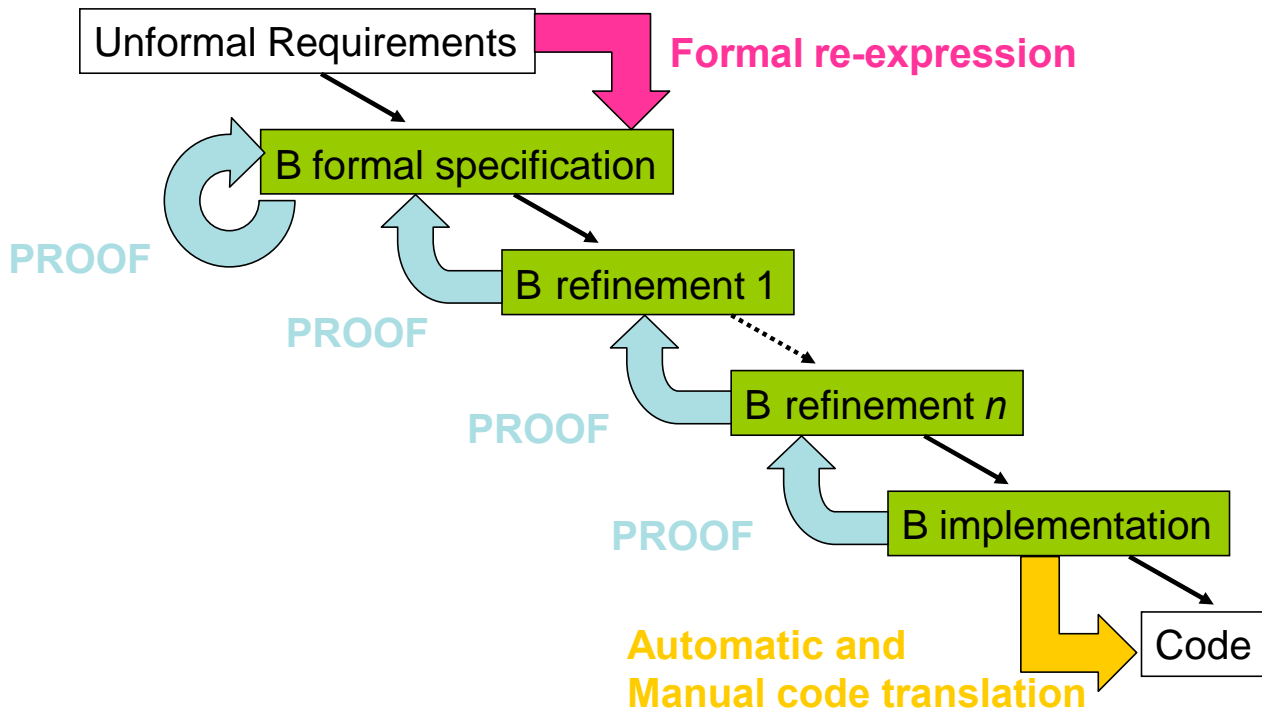


B formal method – examples of safety properties

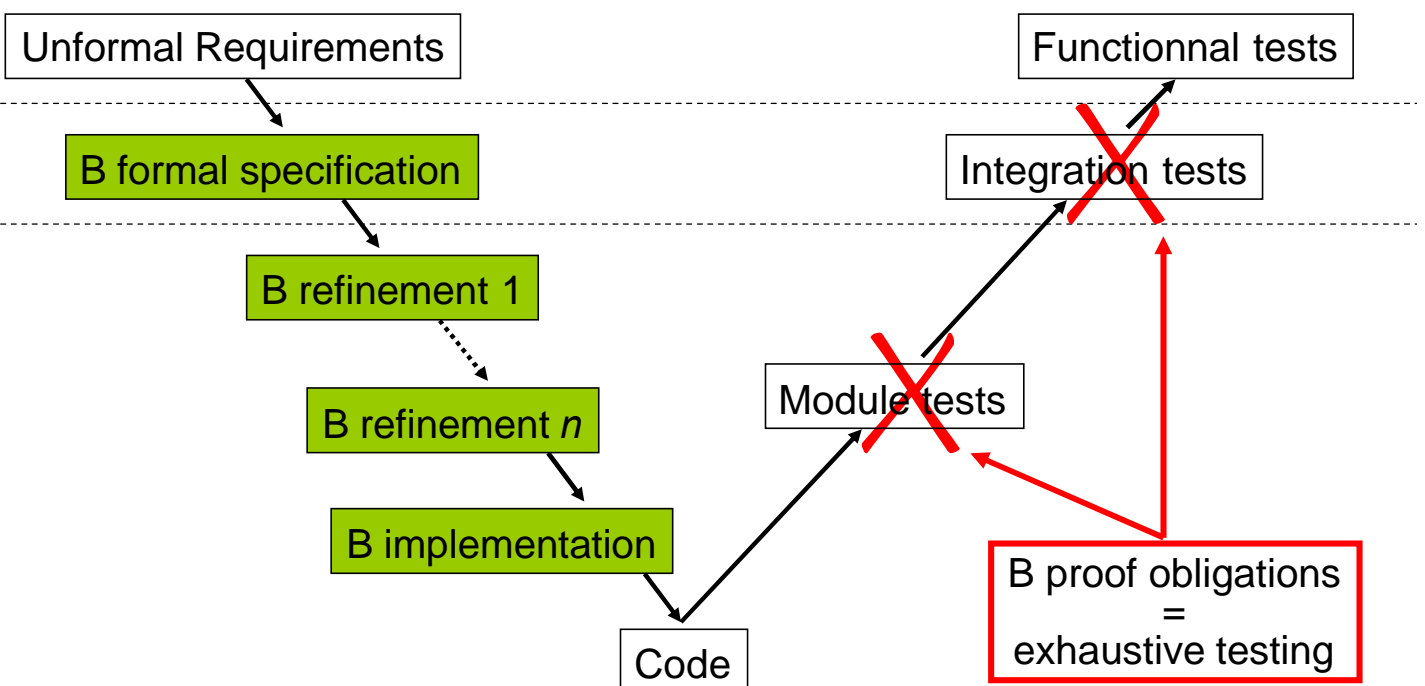
- Only equipped train which is located and in automatic mode can have a target.
- The trains locations computed by the SWE must be correct with the actual trains locations on the line.

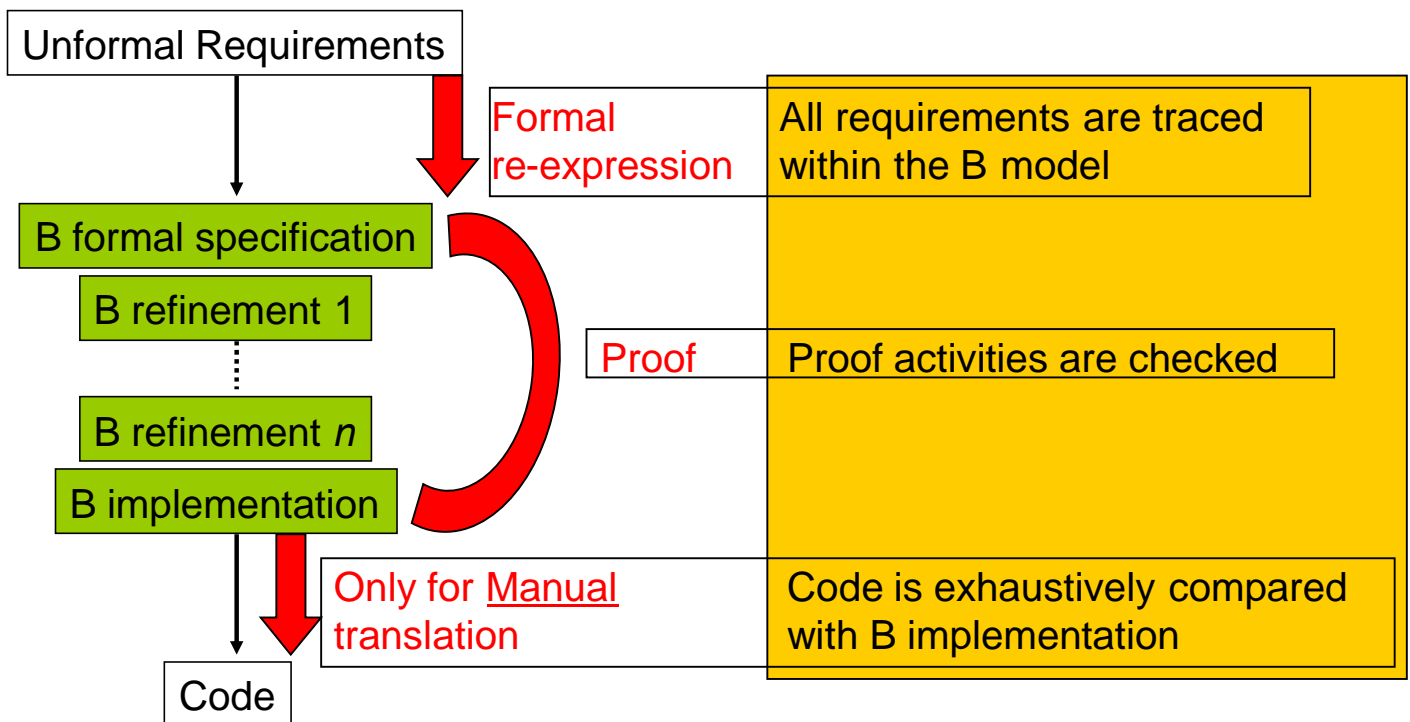


B development process



B verification process



B validation process***B industrialisation***

AtelierB[®] : An industrial tool to specify, refine, implement and prove B models

Statistics about Météor B model

- 1150 B components
- 115 000 lines of B code
- 27 800 Proof Obligations (all proved)
- 86 000 lines of « safe » ADA code

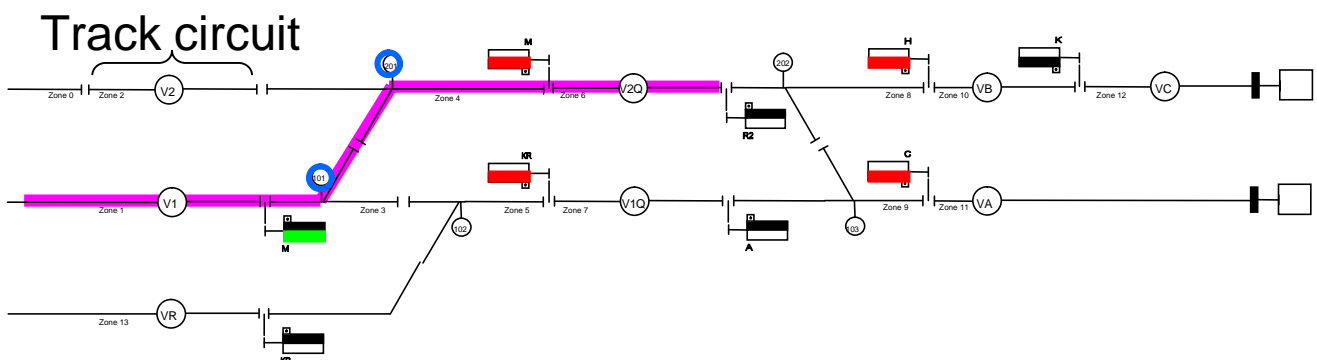
B today in railway industry

Used by two railway leaders : SIEMENS and ALSTOM

Recent projects :

- Canarsie Line (New-York),
- North East Line (Singapour)

Projects size has increased more than twofold

***Interlocking system***

Route V1 to V2Q

Signals M (V2), H, KR, C

Switches 101, 201

Signal M (V1)



Site configuration

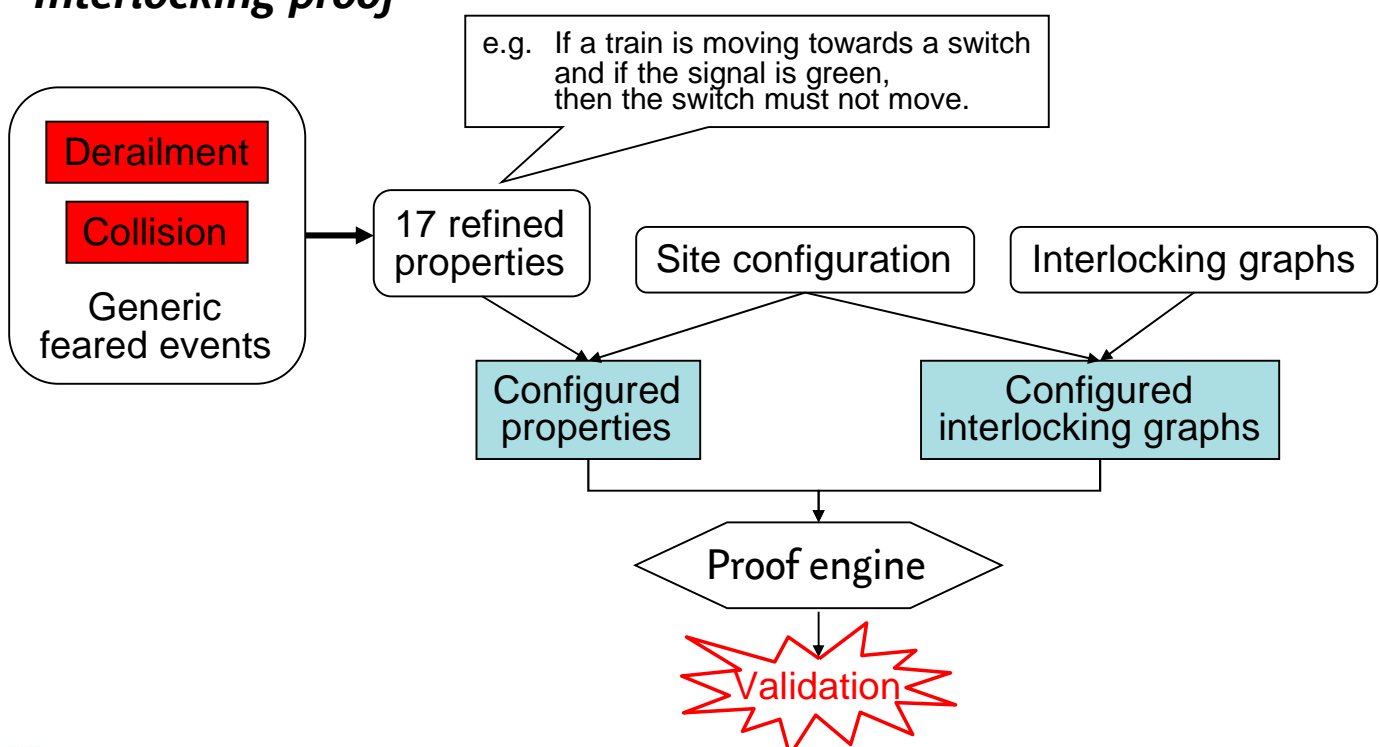
Interlocking validation

Issue: how to be convinced that any combination of generic graphs for any site configuration is safe ?

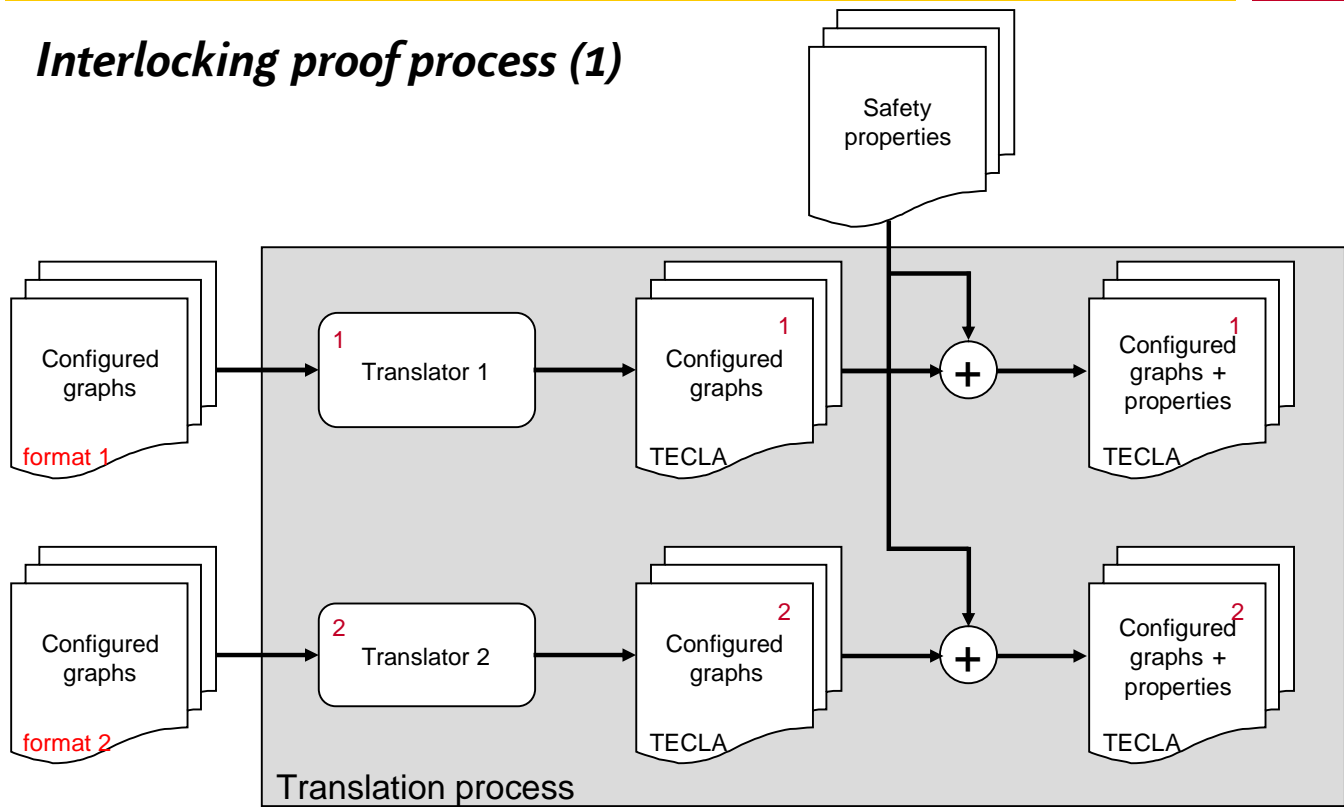
- Heavy testing for both supplier and RATP on site configuration

To reduce test effort for next interlocking sites, formal proof of safety properties has been considered.

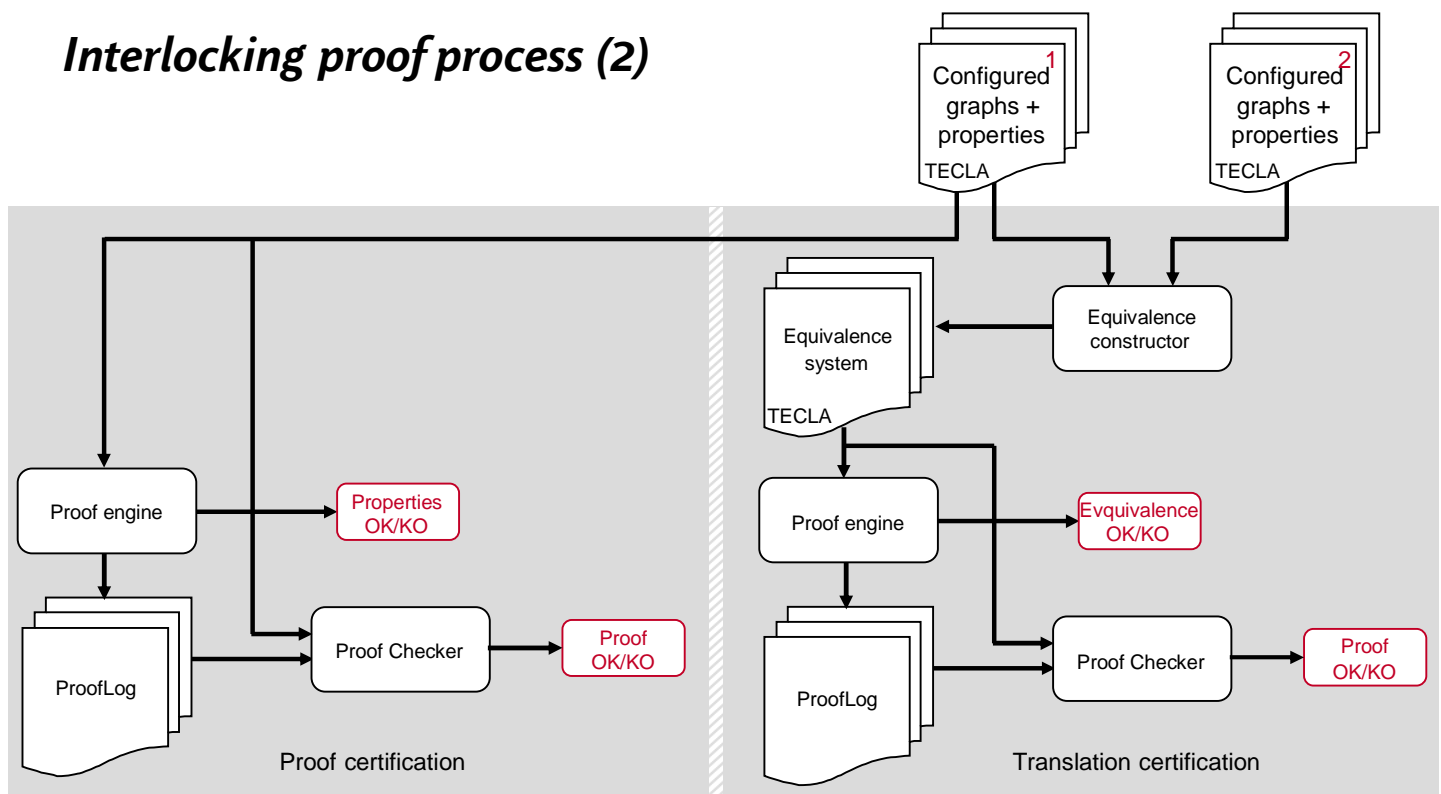
Interlocking proof



Interlocking proof process (1)



Interlocking proof process (2)



Interlocking proof

The proof engine (from Prover Technology) is based on combination of SAT techniques and other automatic proof techniques.

Work in progress

- Feasibility is established
- Complete proof of a real interlocking configuration is expected in a few months

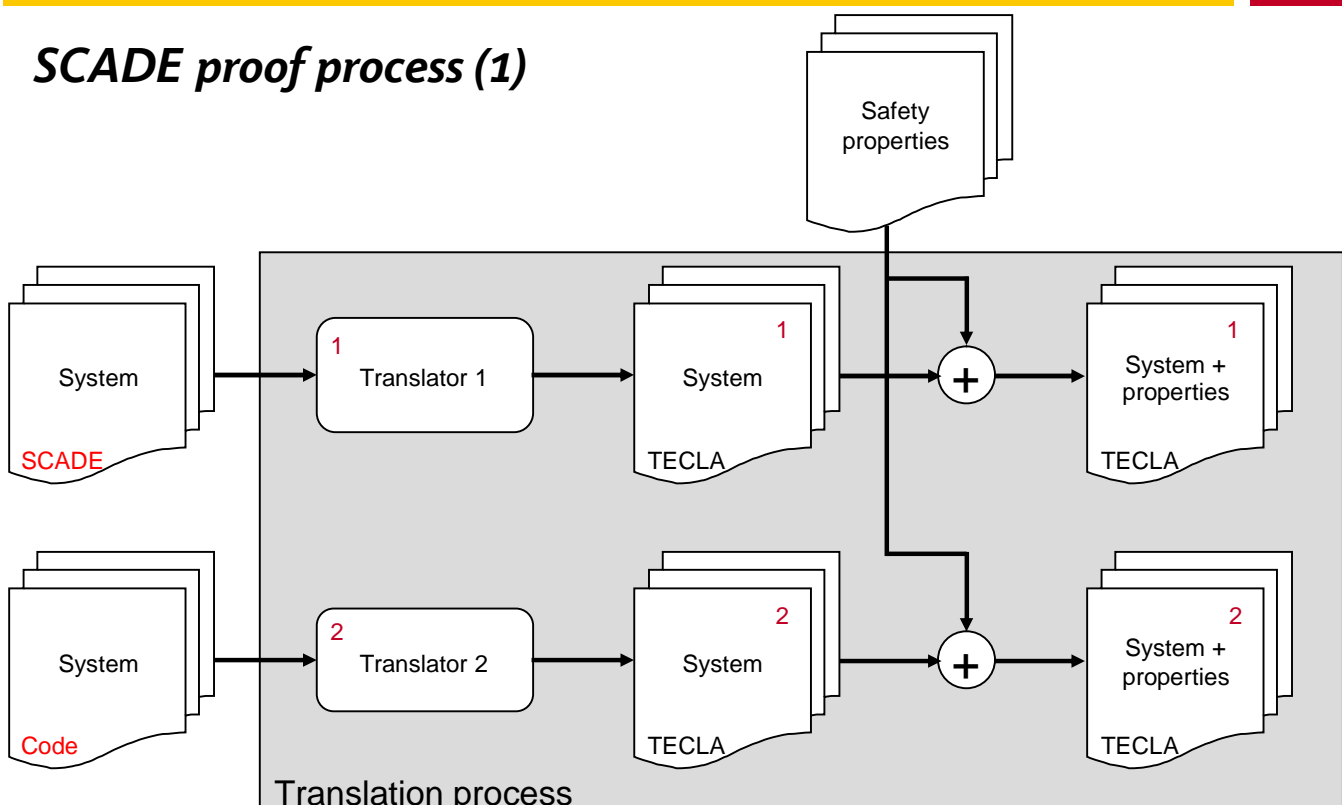
Apparition of SCADE tools in railway industry

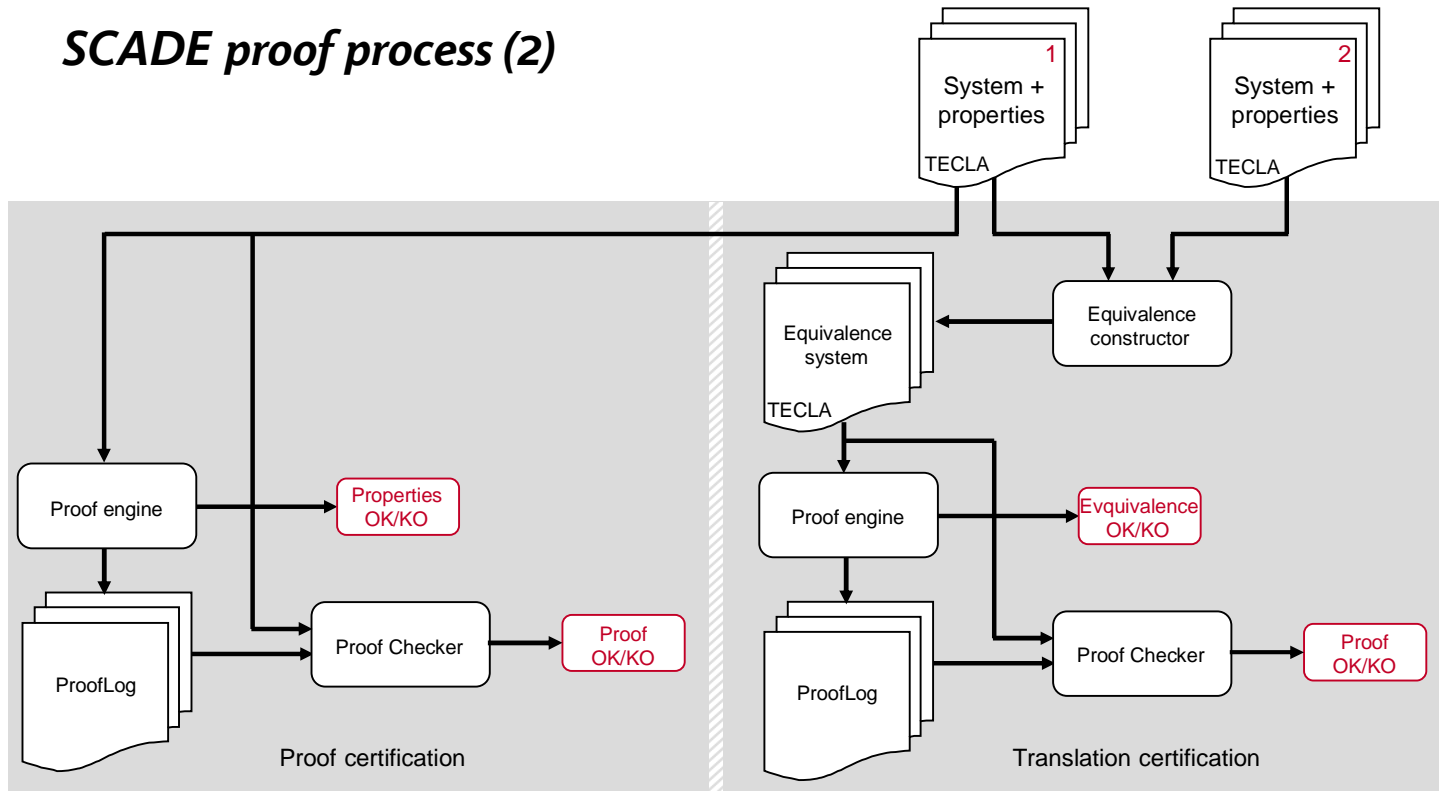
For a few years, SCADE has found favour with railway industry

- fitted for designing command-control systems
- reduces developement cost
- facilitates communication between specialist engineers and software engineers

SCADE brief overview

- based on a declarative synchronous language
Lustre, encapsulated in graphical representation
- Software = variables + equations
- Time is discrete (var_n)_N
clocks, temporal operators (pre, when, ...)
- Equations between inputs and outputs
 $\text{out}_n = \Phi(\text{in}_n, \dots, \text{in}_{n-p}, \text{var}_n, \dots, \text{var}_{n-q})$

SCADE proof process (1)

SCADE proof process (2)**SCADE proof**

Example of safety property :

- Two distinct trains must not cross their movement authority limit

Work in progress








- Feasibility on a real site configuration
- System requirements specification coverage
- Method to complete proof when safety properties are not totally proved



Formal methods ...

- reduce drastically test effort
- provide a high level of quality and safety for software
- are applicable to industrial software projects
- but have to take more into account the practical aspects for using them (cost, competence, ...)

RATP renewal program: software development methods

	OURAGAN CBTC	Manless CBTC
B method	 L3	
Coded processor	 L5 WaySide Equipement	 L1
SCADE	 L5	
Redundant processors	 L3  L13	