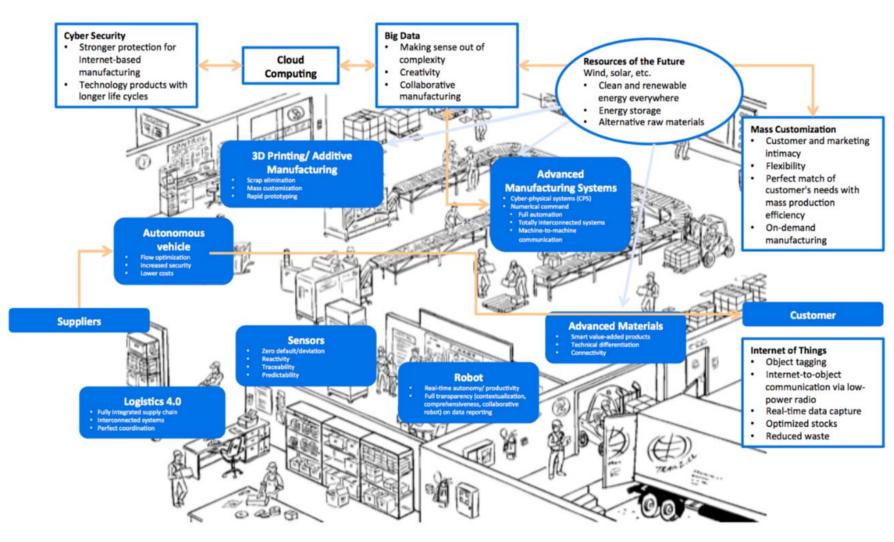
#### Smart Factories and Secure Cloud Storage Services: Vulnerabilities and Countermeasures

Chia-Mu Yu

Yuan Ze University

IFIP WG 10.4 Meeting 25 Jun 2016

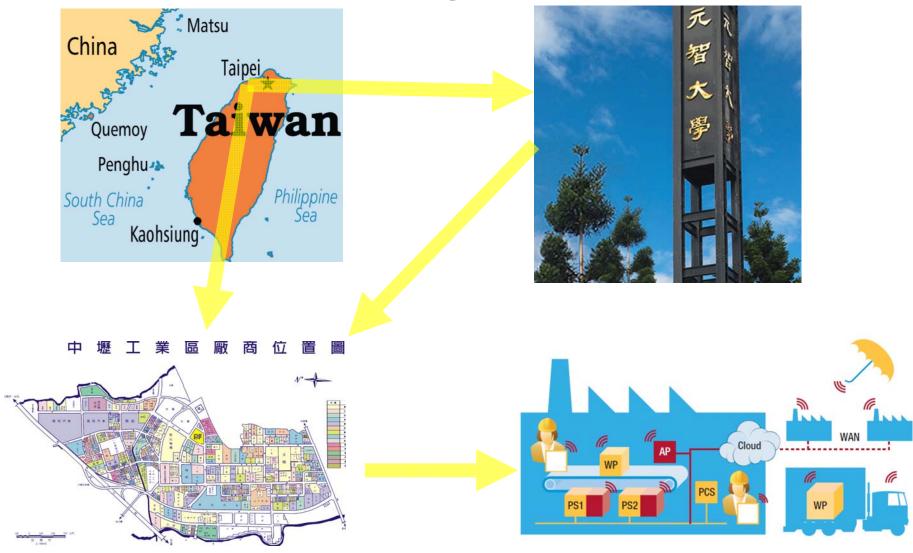
### **Smart Factory**



# Two Cloud Security One IoT Security

will be covered in this presentation

### Motivating Scenario



#### Motivating Scenario

- Restriction
  - My school does not have enough storage
  - Factory owner does not want to release data
  - Factory owner does not have enough storage
- OK... let's try Dropbox (cloud storage)
- Factory owner: Snowden said it's unsafe... Factory owner: I want my data encrypted and efficiency

# Cloud Storage Security

Ah... kind of motivated by smart factory

### Cloud Storage Providers















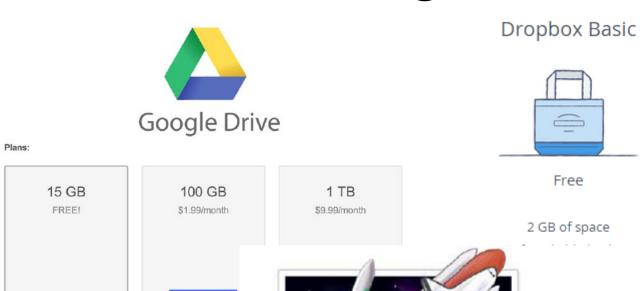








### Cloud Storage Providers



Choose

Current plan

#### Great Space Race!

The Great Space Race has ended! You can see the final results below!

#### Global Leaderboard

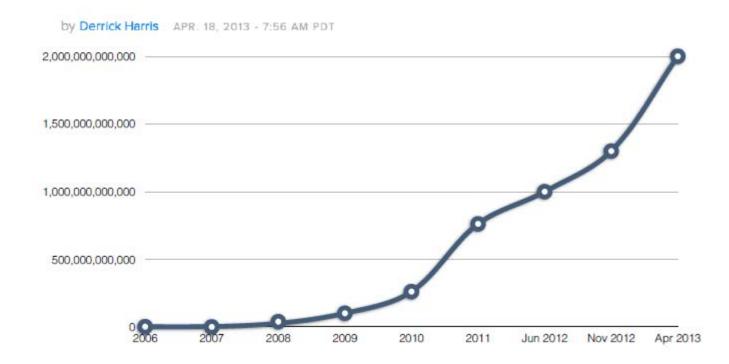
	SCHOOL	NUMBER OF SPACE RACERS	TOTAL POINTS
-11	National University of Singapore	20,532	45,090 points
2	National Talwan University	16,645	40,292 points
3	II Politecnico di Milano	14,425	33,841 points
4	Nanyang Technological University	14,983	33.731 points

# Cloud Storage Providers

Individual Amazon S3 objects can range in size from 1 byte to 5 terabytes. The largest object that can be uploaded in a single PUT is 5 gigabytes. For objects larger than 100 megabytes, customers should consider using the Multipart Upload capability.

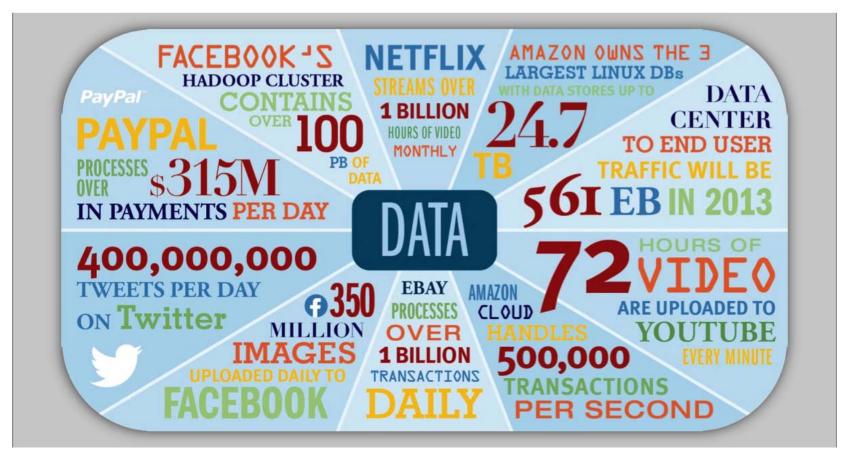
amazon web services / big data / cloud computing

# Amazon S3 goes exponer Amazon S3 FAQs - Amazon Web Services now stores 2 trillion objects



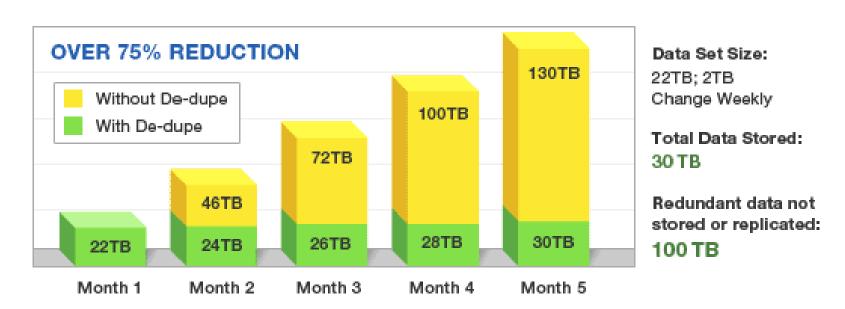
#### Data Deduplication

People keeps uploading stuffs to cloud



#### Data Deduplication

- Data deduplication
  - A way of avoiding storing the same file twice

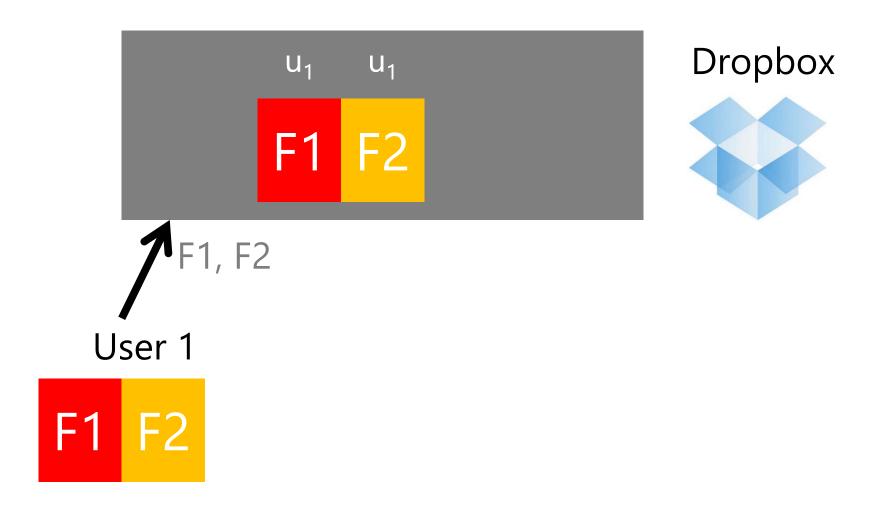


#### Cross-User Server-Side Data Deduplication

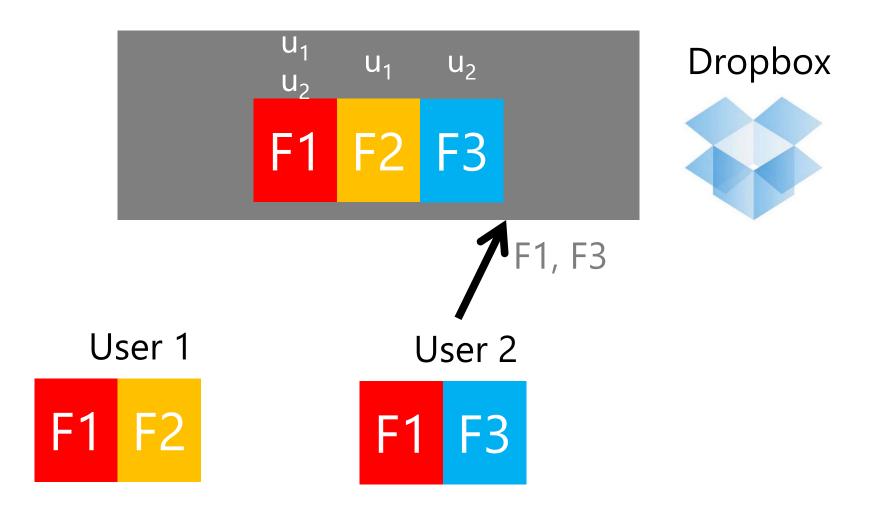


User 1 F1 F2

#### Cross-User Server-Side Data Deduplication

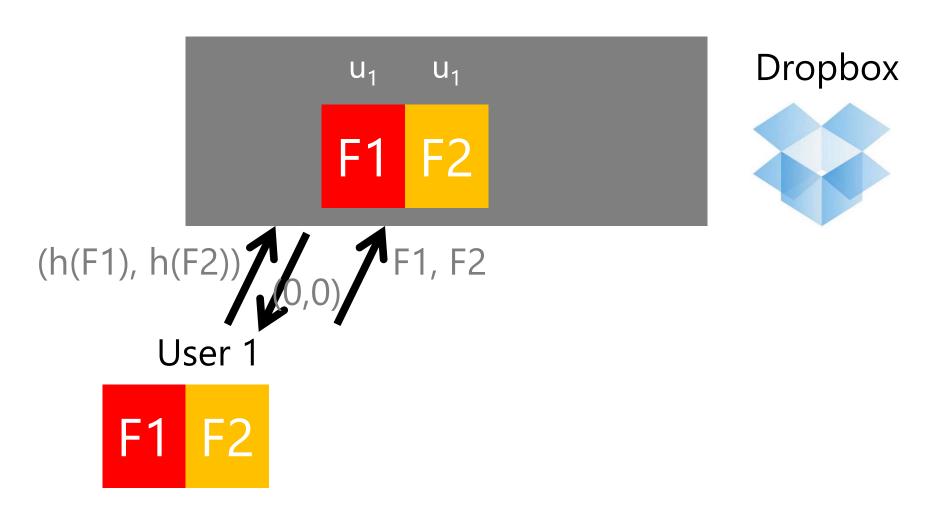


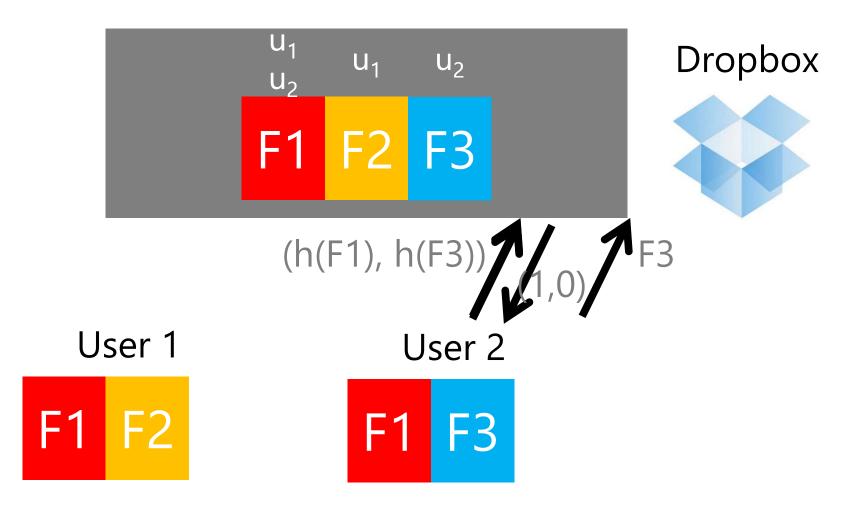
#### Cross-User Server-Side Data Deduplication

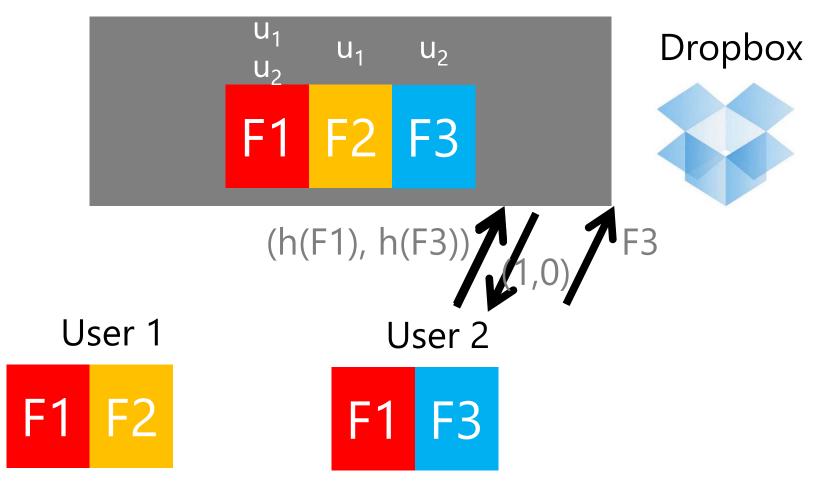




User 1 F1 F2







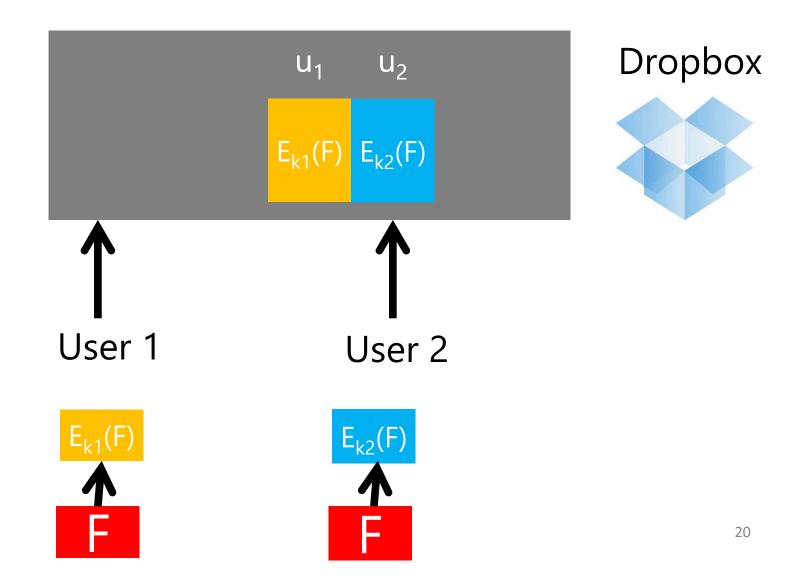
#### Secure Deduplication

- Data could be sensitive
  - Data need to be encrypted before uploaded
  - However, totally destroys deduplication capability



Say DEDUP one more time...

#### **Encryption Meets Deduplication**

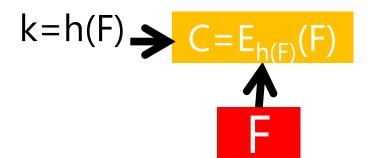


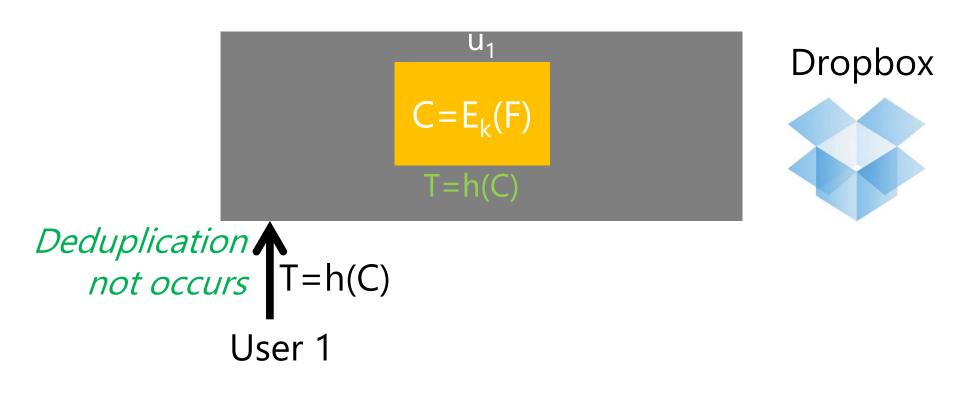


User 1

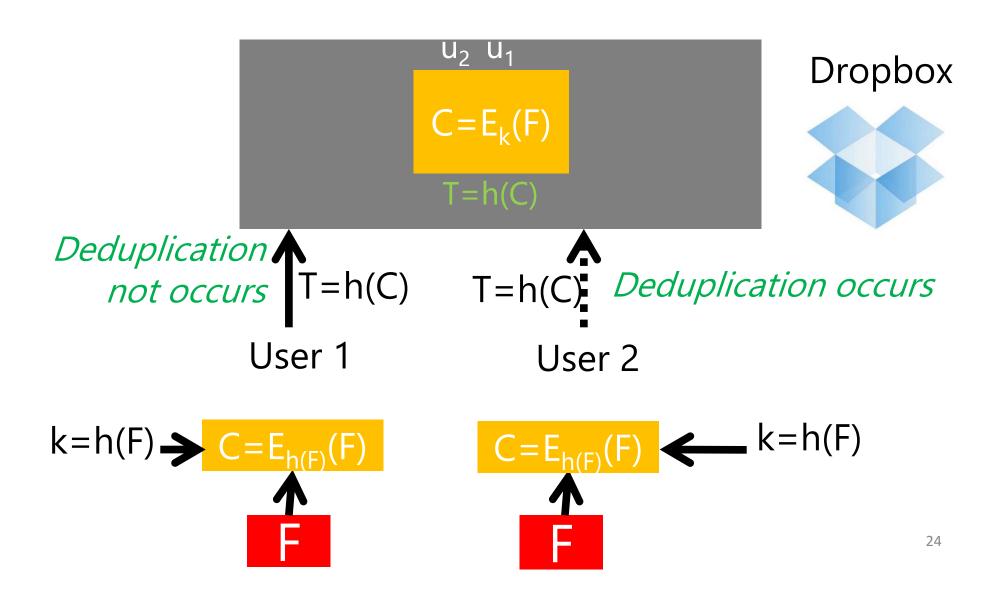


#### User 1





$$k=h(F)$$
  $C=E_{h(F)}(F)$ 



#### Try Every Possible Patterns!

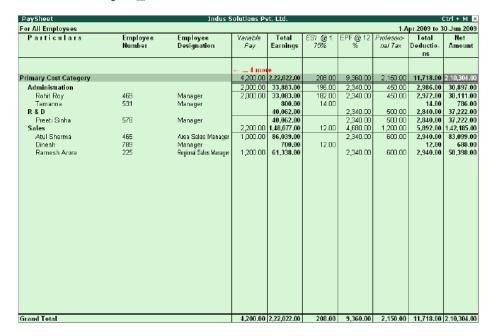
- Convergent Encryption (CE)
  - Good for both data deduplication and privacy

- The weakness
  - -Brute force attack



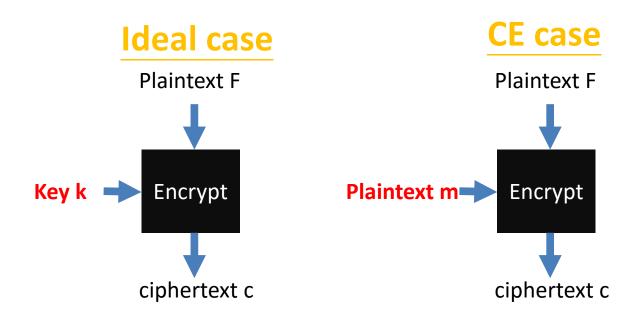
#### Weakness

- File predictability
  - In real life, file content is usually predictable
    - Pay sheet example
    - Chef's secret sauce
    - Engineer's parameter
    - etc



#### Weakness

- Brute force attack
  - MLE is weaker than conventional use of AES
    - Reason is that *CE* is keyless

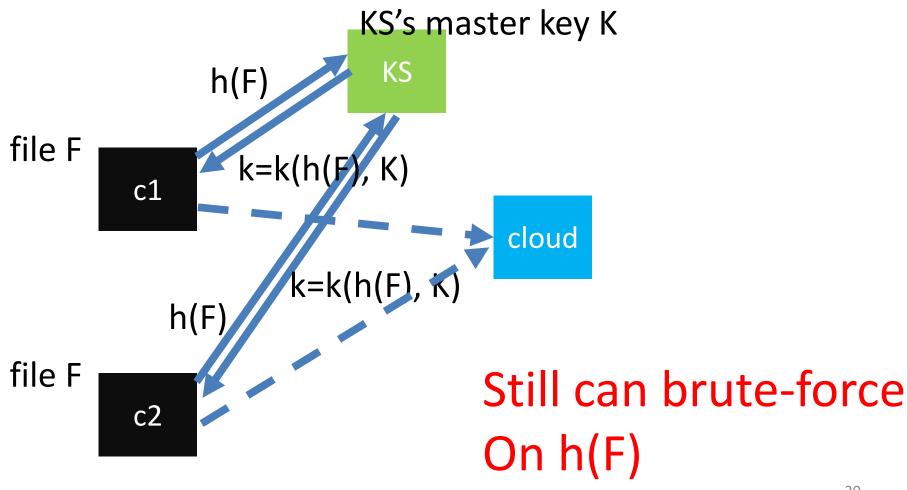


#### Our Requirements

- Data deduplication
- Computation efficiency
- Brute-force resiliency

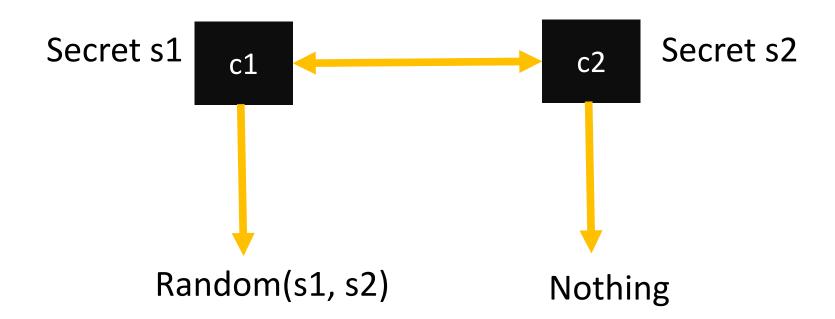
- How to overcome weakness?
  - A new secret
  - Idea is to deploy an additional key server (KS) that is responsible for generating keys for encryption purpose

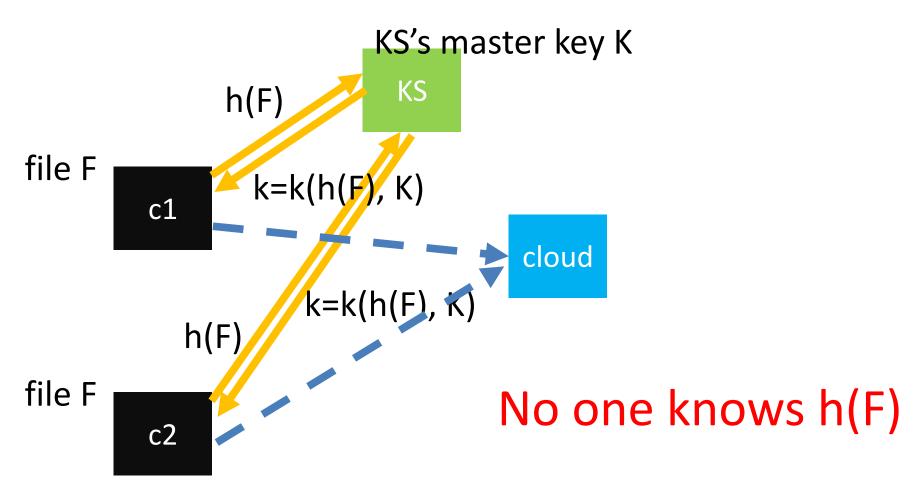
#### Naïve Implementation of DupLESS



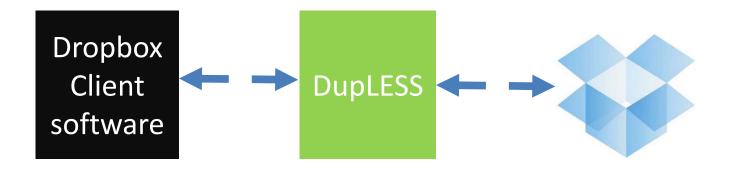
#### **OPRF**

- Oblivious pseudorandom functions
- Kind of blind signature



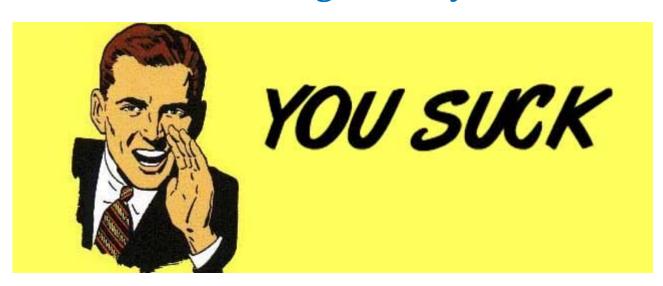


- DupLESS does not need to modify cloud
  - It can be an additional software layer



Factory owner: I'm happy

- DupLESS seems to have no weakness
  - -No
  - -It has no practical use!
  - Who will be in charge of key server?

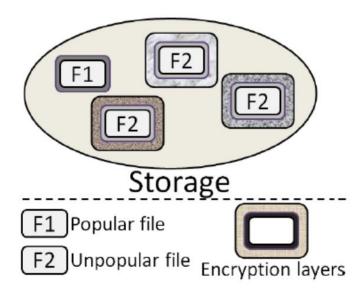


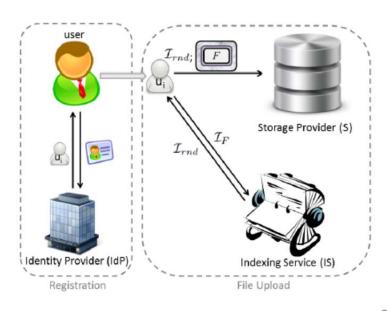
#### SecDep

- DupLESS client always talks to KS, would inefficient in chunk level
  - Upload a file
    - Talk to KS in file level, to get file-level key and check dedup status in cloud
    - If not deduped, talk to KS again in chunk level, to get chunk-level key
- Maintains keys (file/chunk level) by client itself is cumbersome
  - Multiple KS
  - Distribute secret shares of key to KSs

#### Threshold CE

- Dedup according to file popularity
- Each file is encrypted in two layers; the first is, the second is threshold CE



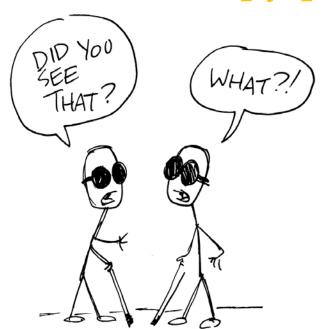


#### No KS Solution?

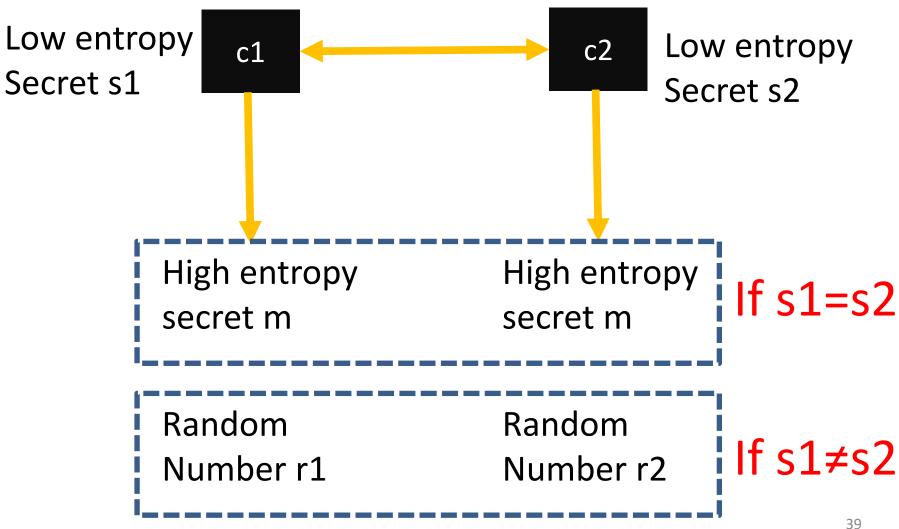
- Where the difficulty from?
  - Cannot send h(f)
    - Brute force attack for low-entropy file f
  - Cannot send E(f)
    - No bandwidth saving
  - Cannot communicate with additional trusted server and communicate via trusted channel
    - Awful assumption

#### **PAKE**

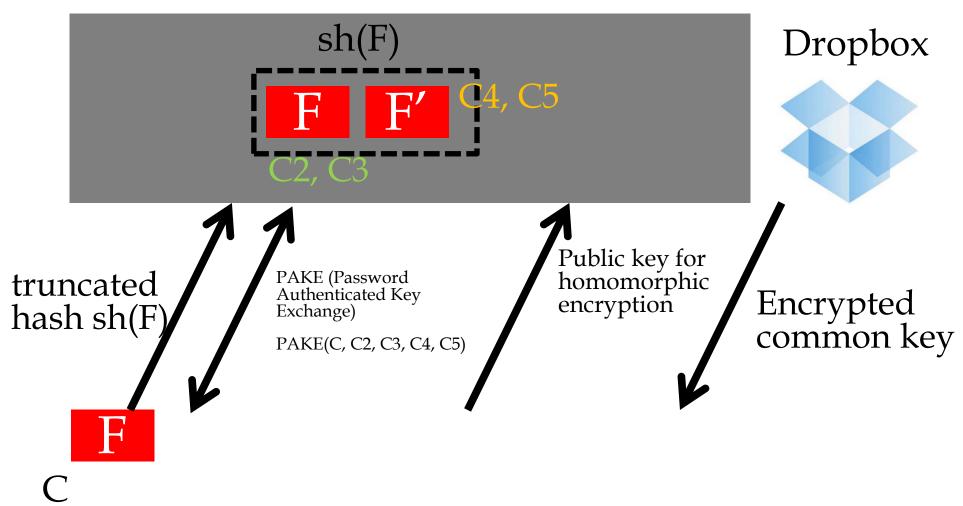
- Password Authenticated Key Exchange
- Enable users to establish a common key based on their low entropy password only



#### PAKE



### PAKE-based Solution



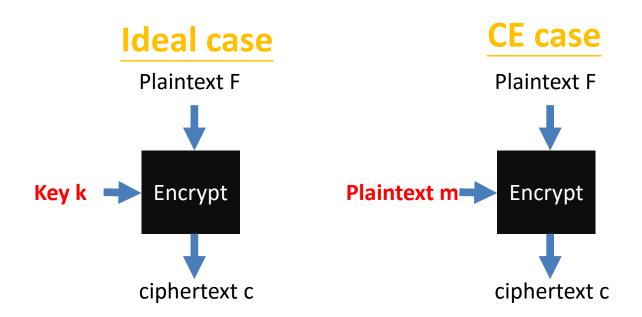
#### PAKE-based Solution

- Two heavyweight weapons
  - PAKE
  - Homomorphic encryption
- Have significant theoretical contribution but still no practical impact



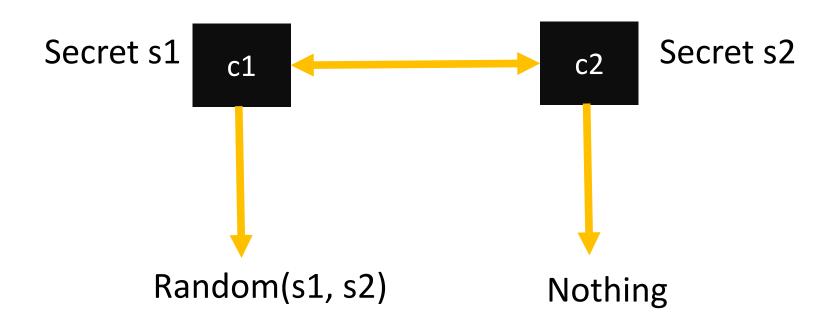
### Rethinking PAKE-based Solution

- In fact, KS is still there; everyone can be KS
  - Essentially, we need an additional secret for brute-force attack

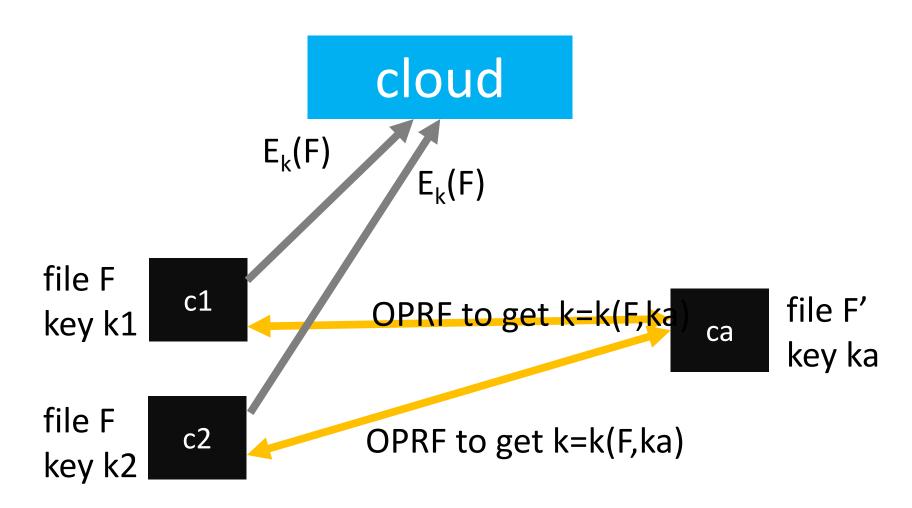


## OPRF, again

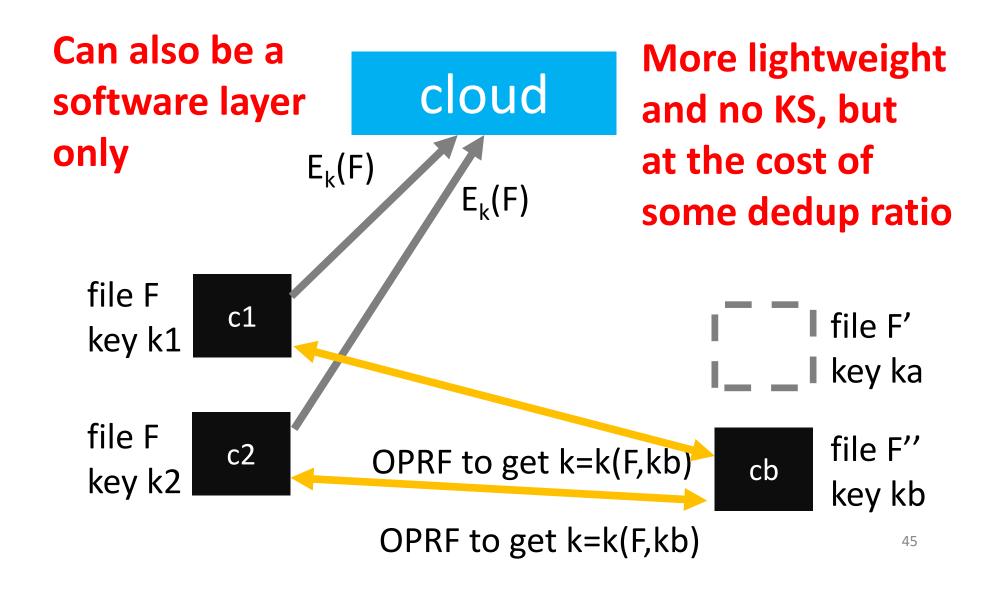
• Combine OPRF and the idea that everyone can be KS



### **OPRF-based Solution**



#### **OPRF-based Solution**

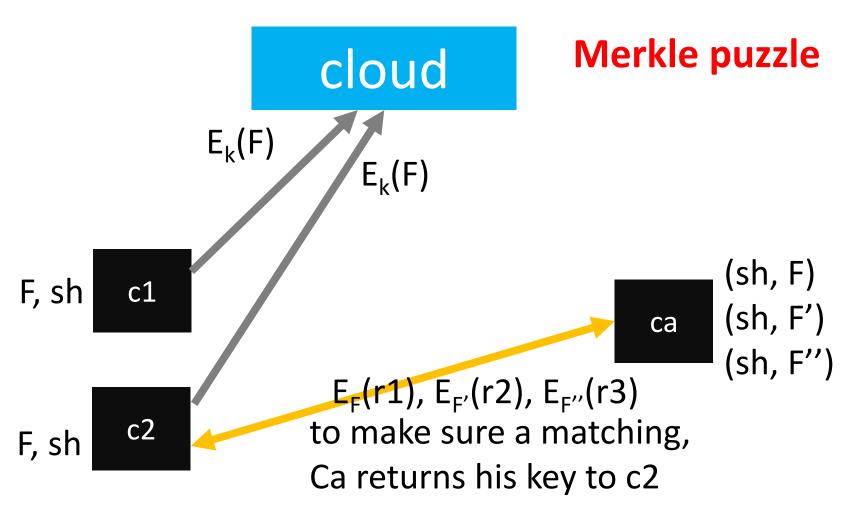


### Symmetric Crypo-based Solution

• Should be the best in terms of performance

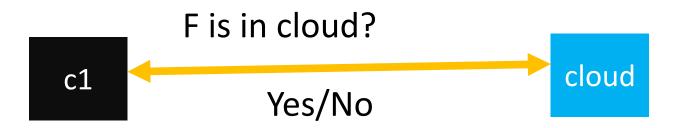
- Take another route, sh(F)
  - E(F), h(F) are not good, OPRF is heavyweight

### Symmetric Crypo-based Solution



## Motivating Scenario

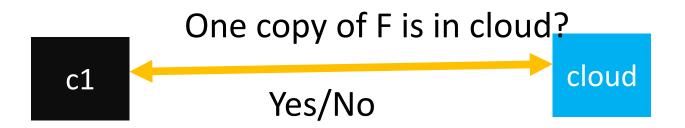
- Factory owner: dedup leaks my secret
- Me: why?
- Factory owner: cloud always returns dedup result!



### Threshold = 1

 Original deduplication assumes threshold=1

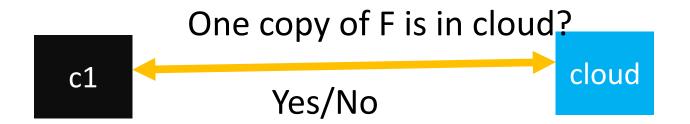
• Easy for attacker to know the file existence status



#### Random Threshold

• Each file x is associated with a random threshold tx

 tx too large, no dedup tx too small, no security



• First attempt: randomize the hash response

Chunk existence	Hash response
0	1
1	0/1

• 0-response indicates chunk existence



• Second attempt: client uploads two chunks at once

Chunk 1	Chunk 2	response
0	0	2
0	1	1
1	0	1
1	1	1

- Upload **c1⊕c2** to cloud
- Seem to work?
- Fix a chunk not in cloud, infer chunk 2 existence

• Third attempt: each result has a time limit

Chunk 1	Chunk 2	response
0	0	2 (t)
0	1	1 (t)
1	0	1 (t)
1	1	1 (t)

• Many accounts query cloud within a short time period

 Fourth attempt: client cannot do the query but does not upload the chunk

Chunk 1	Chunk 2	response
0	0	2
0	1	1
1	0	1
1	1	1

Many accounts will be used by attacker

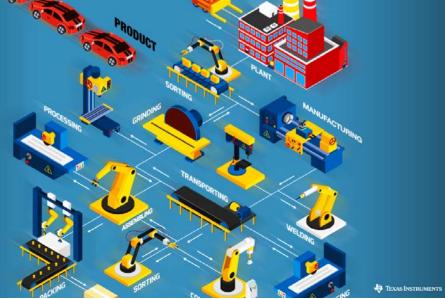
Observation: in any case, at least c1⊕c2
needs to be uploaded

Chunk 1	Chunk 2	response
0	0	2
0	1	1
1	0	1
1	1	1

• Force client to send the query with the form of  $(h(c1), h(c2), c1 \oplus c2)$ 

# Motivating Scenario





# Internet of Things Rule Checking

# Augmented Collective Beings

• There are a lot of devices interacting with each other and with users, who are usually not IT professionals.



### Cross-device Dependencies

- Explicit dependencies
   If power usage is higher than 50, turn off air conditioner
- Implicit dependencies
  - Via context, like temperature, location, human
     If air conditioner is turned off, temperature
     increases



## Multi-stage Attacks

 Emerging threats via exploiting explicit/implicit dependencies to access higher-value targets

e.g., burglar wishing to break in can first turn off smart plug, which disconnects the air conditioner, which increases the temperature, which then triggers the window to open.



## Objective

• Given a bunch of dependency rules, check whether several security and safety constraints are violated

### Related work in firewall checking

- [1] checks anomalies that could exist in a single- or multi-firewall environment
- [1]: rules in sequence
- The execution order of firewall rules is fixed with respect to each packet.
   However, every rules operate in parallel in IoT

### Related work in SDN

• [2] is a layer between SDN controller and network devices that checks for network-wide invariant violations dynamically as each forwarding rule is inserted.

• The search space in SDN is fixed to the space of IP headers.

However, in IoT, the search space changes when devices join or leave.

### Related works in IoT

- Most works focus on checking the existence of **conflicts**, which means that multiple rules try to use one or more sensors or actuators at the same time, which cause different effects on the environment
  - Conflicts between rules: [3], [4], [5]
  - Conflicts between users: [6]
- May not be applied directly
  - Global constraints may not be converted to rules
  - The conflicts between pairs may be too strict

<sup>[3]</sup> Policy conflicts in home automation @ Computer Networks: The International Journal of Computer and Telecommunications Networking 2013

<sup>[4]</sup> DepSys: Dependency Aware Integration of Cyber-Physical Systems for Smart Homes @ ICCPS '14: ACM/IEEE 5th International Conference on Cyber-Physical Systems

<sup>[5]</sup> An Application Conflict Detection and Resolution System for Smart Homes @ 2015 IEEE/ACM 1st International Workshop on Software Engineering for Smart Cyber-Physical Systems

<sup>[6]</sup> Conflict detection and resolution in home and building automation systems: a literature review @ Journal of Ambient Intelligence and Humanized Computing October 2014

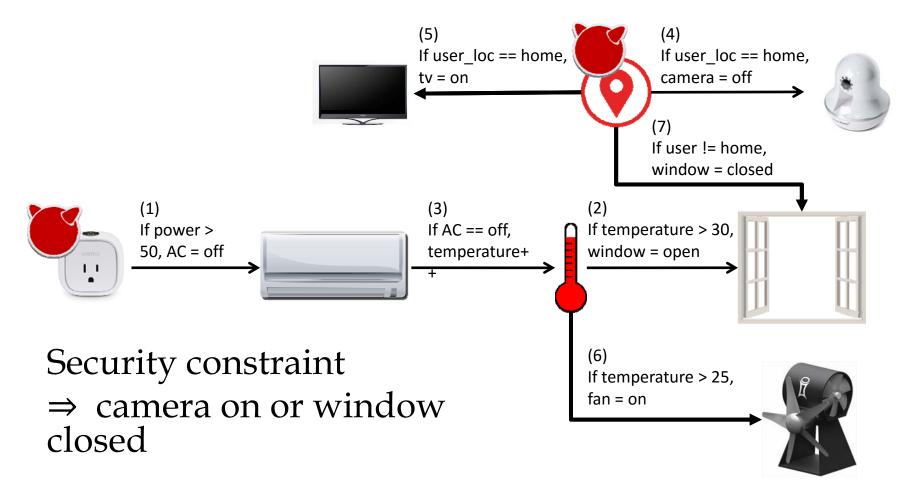
### Related works in IoT

- [7] first considers the security challenges of cross-device dependencies in IoT
- [8] is mostly related

Build a safety-centric programming platform for connected devices in IoT environments. However, the solution they proposed is not fast enough

[7] Handling a trillion (unfixable) flaws on a billion devices: Rethinking network security for the Internet-of-Things @ HotNets-XIV Proceedings of the 14th ACM Workshop on Hot Topics in Networks Article No. 5, 2015 [8] SIFT: Building an Internet of Safe Things @ IPSN '15 Proceedings of the 14th International Conference on Information Processing in Sensor Networks, 2015

### Small Dataset

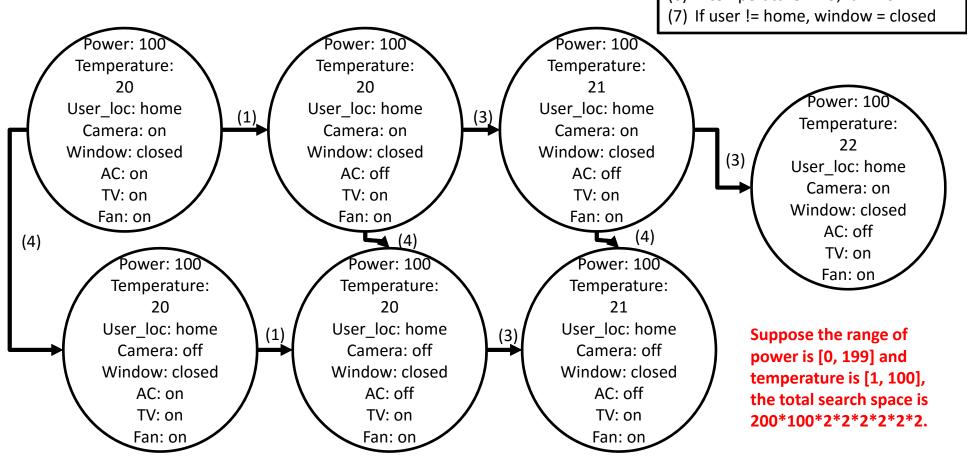


## First try

- The "if-this-then-that" clause is similar to the "implication" in logic
  - ⇒ try to model the rules in the form of propositional logic
- However, the concept of "state" is absent in simple logic
  - ⇒ the situation in which the temperature or the power usage increases cannot be modelled.
- Thus, use state machine to model the rules' effect on environments

# Finite State Machine 10 If power > 50, AC = off

- (2) If temperature > 30, window = open
- (3) If AC == off, temperature++
- (4) If user\_loc == home, camera = off
- (5) If user\_loc == home, tv = on
- (6) If temperature > 25, fan = on



### Conclusion

- Three security issues
  - Two for cloud
  - One for IoT
  - IoT security is more related to smart factory in a straightforward way





